

1 EAs and SAT

In the Boolean Satisfiability Problem (SAT), one has a Boolean expression composed of conjunctions and disjunctions of literals. A literal is a Boolean variable or the negation of a variable. The variables v_i can be true or false. An example (in prefix notation) is $(\wedge (\vee v_1 \neg v_2 v_3) (\vee \neg v_1 \neg v_4 \neg v_5))$. One class of these problems is called “3-SAT”, which has the form given in the previous sentence. Expressions have the form $(\wedge clause_1 clause_2 \dots clause_M)$. Each clause has the form $(\vee . . .)$ where each “dot” is a literal. The goal is to find any assignment to the variables such that the expression evaluates to true. There are N variables and M clauses.

1(a) EAs for 3-SAT

Assume the problems are 3-SAT in nature.

If we solve this with a genetic algorithm, an individual can simply be a bit string of length N . Bit i is 1 if v_i is true, or bit i is 0 if v_i is false. Note that mutation and crossover will work fine. Define a fitness function that evaluates better individuals as having higher fitness, with the maximum fitness being equal to 1.0 (this would be a solution that satisfies the whole expression). Give me a function that will work with general 3-SAT problems of N variables and M clauses.

1(b) EAs for SAT

Now assume that the Boolean expressions have arbitrary form. For example, we could have $(\wedge \neg(\vee v_1 v_2) v_4 (\vee \neg(\wedge v_2 v_3) v_5 (\wedge v_5 \neg v_6)))$. In other words, the expression is an arbitrary combination of ANDs, ORs, NOTs and variables. Note that this time whole expressions can be negated. Again assume an individual is a bit string of length N , with 1 = true and 0 = false.

Define a fitness function that evaluates better individuals as having higher fitness, with the maximum fitness being equal to 1.0 (this would be a solution that satisfies the whole expression). Give me a function that will work with general SAT problems, not just the trivial example above.