Dominion: An Architecture-driven Approach to Generating Efficient Constraint Solvers

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***Abstract*—Constraints are used to solve combinatorial problems in a variety of industrial and academic disciplines. However most constraint solvers are designed to be general and monolithic, leading to problems with efficiency, scalability and extensibility. We propose a novel, architecture-driven constraint solver generation framework called Dominion to tackle these issues. For any given problem, Dominion generates a lean and efficient solver tailored to that problem. In this paper, we outline the Dominion approach and its implications for software architecture specification of the solver.**

I. Introduction

Constraints are a natural and powerful means of representing and reasoning about combinatorial problems. Constraint solving provides a mechanism for finding solutions to such problems automatically. Its simplicity and generality are fundamental to its successful application in a wide variety of disciplines such as scheduling, industrial design, aviation, banking, combinatorial mathematics and the petrochemical and steel industries [1].

Constraint solving of a combinatorial problem proceeds in two phases. First, the problem is modelled as a set of decision variables and a set of constraints on those variables that a solution must satisfy. A decision variable represents a choice that must be made in order to solve the problem. The domain of potential values associated with each decision variable corresponds to the options for that choice. The second phase consists of using a constraint solver to find solutions to the model: assignments of values to decision variables satisfying all constraints. Constraint solvers typically employ a systematic backtracking search through the space of partial assignments in order to find solutions.

Most current constraint solvers, such as Minion [2], are constructed to be as general as possible. They are monolithic in design, accepting a broad range of models. While this generality is convenient, it leads to a complex internal architecture, resulting in significant overheads and inhibiting efficiency, scalability and extensibility. Another drawback is that current solvers perform little analysis of an input model, so the features of an individual model cannot be exploited to produce a more efficient solving process. To mitigate these drawbacks, constraint solvers often allow manual tuning of the solving process. However, this requires considerable expertise, preventing the widespread adoption of constraint solving.

The main aim of the work introduced in this paper is to improve the scalability of constraint technology, while simultaneously removing its reliance on manual tuning by an expert. We propose a novel, elegant means to achieve this aim: a constraint solver generator framework called *Dominion* which, for a given problem, produces a solver tailored to that problem. There are two key benefits to this approach:

1) it enables fine-grained optimisations not possible for a general solver, allowing the solution of much larger, more difficult problems, and
2) it enables the utilisation of many techniques in the literature that, although effective in a limited number of cases, are not suitable for general use, leading to more powerful solvers.

The Minion solver, mentioned earlier, allows some specialisation of components. There are 7 variable types, and each constraint is compiled for two inputs which can each be different variable types, therefore currently each constraint is compiled 49 times. Adding one extra option to variables (doubling the number of variable types) increases compilation time fourfold. In contrast, Dominion compiles exactly the required variables and constraints for each problem.

The generation process in Dominion is driven by the software architecture of the target solver. Software architecture provides a high-level model of the structure and behaviour of a system in terms of its constituent elements and their interactions as well as conditions that have to hold among the elements [3], [4]. Desired properties of a system can be checked at the architectural level as well as against an implementation. Thus the architecture forms a useful basis and guide for the implementation and evolution of systems.

A number of Architecture Description Languages (ADLs) have been developed to capture these features [5]. We use a powerful, general-purpose ADL called Grasp in this work. Grasp provides support for commonly used generic architectural primitives such as layers, components, connectors, interfaces, properties and wiring as well as the ability to define user-defined constructs for particular domains. It allows the specification of templates which can be instantiated to produce the required architectural elements. Most importantly, unlike many existing ADLs, Grasp has been designed to capture sophisticated dependencies among architectural elements. These features make Grasp ideally suited to represent the architectures of customised constraint solvers in Dominion.

In this paper, we present an outline of our approach and focus on its implications for software architecture specification.

II. RELATED WORK

One of the earliest examples of systems that attempt to generate constraint solvers tailored to a specific problem is the MULTI-TAC system [6], which configures and compiles a constraint solver for a specific set of problems. It is written in LISP and performs ad-hoc customisation of a base constraint solver limited to a few characteristics.

KIDS [7] is a more general system that also uses LISP to synthesise efficient algorithms from an initial specification. The approach is knowledge-based, i.e. the user supplies the knowledge required to generate an efficient algorithm for the specific problem. Refinements are limited to a number of generic transformation operations. Our approach is more general and, crucially, relies on almost no background knowledge.

There are other constraint solver systems that perform code generation or modification, but to the best of our knowledge no previous approach to formally specifying the architecture of specialised constraint solvers exists.

III. AN EXAMPLE

The following listing shows a very short example of a constraint problem expressed in the Dominion Input Language (DIL) [11]. It will be used in the remainder of the paper to illustrate architectural support for the Dominion approach.

```
language Dominion 0.1

given b : int {1..}
given s : int {1..}

find w : int {2}
find x : int {0..1}
find y : int {0..b}
dim z[s] : int
find z[..] : int {-2,0,1}

such that
conA sum([x,y], w)
conB sum(z[..], w)
```

Listing 1: A Constraint Problem

The basic format of a problem is to give a list of parameters (`given`), decision variables (`find`), their domains and finally a list of constraints the variables must satisfy (`such that`). The `dim` statement denotes that z is a one dimensional matrix of size s. Given the values $b = 2$ and $s = 2$ for the parameters, one solution to this problem is $w = 2, x = 1, y = 1, z[0] = 1, z[1] = 1$. Different stages of the Dominion solver generator process will be illustrated using this example, beginning with a description of issues that arise in modelling constraint solvers.

IV. ISSUES IN MODELLING CONSTRAINT SOLVERS

A constraint solver is typically built from a number of factories, which create the various parts of the solver and the connections between them.

Two of the most important parts of a constraint solver are the components which represent the (decision) *variables* and *constraints*. Variable components maintain a set of values the variable might take in a solution (the *domain*). The constraints query variable domains, and remove values that can take part in no solution.

To illustrate some of the key problems in architectural modelling of constraint solvers, we turn to the Minion solver, introduced earlier. In Minion there are several variable factories, producing variables with different features and performance characteristics. For example, there is a factory for variables with initial domain $\{0, 1\}$. Some other variable factories require the initial domain to be an unbroken range. Some implementations permit values to be removed from the middle of the domain, while others only allow changes at the upper and lower bounds. Constraints place callbacks on variables to indicate when they should be informed of changes. Some constraints change the callbacks during search (*dynamic* callbacks) while others do not. It would be desirable to specialise variables on whether they allow dynamic callbacks. Minion does not do so because each specialisation adds to its compile time.

Minion may query variable domains hundreds of millions of times per second, therefore tiny changes to variable implementations can lead to large changes in solver performance.

Constraints also have restrictions upon them. As an example, consider the constraint $\sum X = y$, where X is a vector of variables. The components below, among others, have all been implemented in Minion for this constraint.

1) `GACSum`: An exponential time algorithm which reduces domains as much as possible. Requires all variables to allow arbitrary domain removals.
2) `BoundSum`: A generic polynomial time algorithm.
3) `BoolSum`: Requires all variables in X to have domain $\{0, 1\}$.
4) `BoolSumConst`: Requires all X to have domain $\{0, 1\}$, and support dynamic callbacks. Also the domain of y must be a single value.

Some requirements, such as dynamic callbacks or arbitrary domain removals, are restrictions on the type of variable used. Other requirements, for example in `BoolSum` that variables have domain $\{0, 1\}$, are not restrictions on the type, but on the variable. While there are variable factories which only produce variables with domain $\{0, 1\}$, any variable with domain $\{0, 1\}$ will be accepted by `BoolSum`.

These options are wired into Minion in different ways. `GACSum` can be explicitly chosen by users. The others are chosen using simple heuristics, which are known to be incorrect in some cases. These heuristics may pick an implementation which requires dynamic callbacks, so all variables are required to support them.

Thus, an important requirement of an architecture description of solvers is the ability to associate properties with components and specify and check dependencies among them.

V. ARCHITECTURE SUPPORT FOR DOMINION IN GRASP

Software architecture is the main driver of solver generation in Dominion and contains much of the vital information required for the process. This has implications for the expressive power of the chosen ADL. In addition to customary details of components and connections, the Dominion approach requires further support from the ADL in order to automate the

process of generating an optimal architecture from the problem component, and the solver code from the architecture.

Grasp is a general purpose textual ADL designed to capture architecture rationale as well as structural and behavioural aspects. Supported architectural primitives include layers, components, connectors, templates, interfaces, rationale, wiring, properties and check clauses. Templates are abstractions for architectural elements and can be used to create instances with common behaviour. The required and provided interfaces allow tools to match the functionality of linked components. Properties are characteristics (or restrictions of functionality) associated with architectural elements in the form of name-value pairs, while check clauses allow properties of parameters and linked elements to be checked for compatibility. Grasp also provides a generic annotation mechanism to associate meta-data with architectural elements. In Dominion, this mechanism is used to specify the locations and file names of corresponding implementations for components.

Listing 2 shows a sample template specification in Grasp:

```
@Dominion(Location = "/lib/src/", Filename = "bool_var.hpp")
template BoolVariableFactory() {
  provides IPropVariable;
  provides IRemoveFromDomain;
  requires IMemoryManager mmanager;
  check mmanager->properties() subsetof
    [(MemoryChanges, 'Single')];
  property domainSize = 2;
}
```

Listing 2: A Template in Grasp

This template is for a factory that creates boolean decision variables, which implement the interfaces `IPropVariable` and `IRemoveFromDomain`, and require to be connected to a component that implements `IMemoryManager`. The `check` statement states that the memory manager is only allowed to have the property (MemoryChanges, 'Single'). The domain size of boolean variables is 2.

The Grasp toolset currently consists of a parser and a checker. Further tools for visualising and validating architectures as well as performing traceability analysis are planned for the language.

VI. The Generative Process

The process of generating an efficient and lean solver for a given problem is driven by its software architecture and contains the following steps as illustrated in Figure 1:

- Problem component generation,
- Architecture generation and analysis,
- Solver generation, and
- Execution monitoring

The overall process can be considered as a control loop with the problem specification initiating feedforward control and data on the execution of the solver leading to feedback control. C++ is used as the implementation language for performance, modularity and backwards compatibility reasons. The solver is developed using component-based software engineering practices. The component library, which is used throughout the solver generation process, is introduced first. The first 3

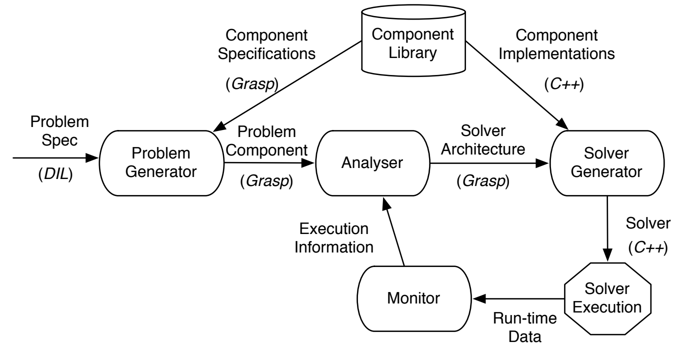

Fig. 1: The Generative Process for Creating Solvers

steps shown in Figure 1 are discussed in the later subsections. Execution monitoring will be added as part of future work.

A. *Component Library*

Components form the building blocks of Dominion constraint solvers. A large number of reusable components are kept in a library to aid the different stages of the generation process. The Dominion component library consists of two parts: the specification of components as elements at the architectural level and the corresponding implementation of these components at the code level. Listing 2 is an example of the architecture level specification of a component template. Each component is implemented as a C++ class, appropriately parameterised to allow customisation as required. Each class is stored in a separate file, increasing flexibility of use and modularity. This practice also allows the generation of the final solver to be automated in a relatively straight-forward manner.

B. *Problem Component Generation*

The generative process is initiated by a constraint problem being supplied to the problem component generator tool. This tool parses the DIL specification and defines a new component in Grasp capturing the essence of the problem such as interfaces that should be supported by variable and constraint components, and any restrictions, such as domain compatibility, among them. Each problem component is unique to the constraint problem to be solved and thus such components are not stored in the component library.

For the example from Listing 1, the problem component generator will produce the specification shown in Listing 3.

The problem component will require to be linked to appropriate variable and constraint components. The check clauses of this component can be divided into two parts. The first places restrictions on the variables based on the domains they will have. For some variables, such as w, the exact domain and size of domain is known. For others such as y, the type of domain is known but not its size. Thus the choice of component for the variables is constrained to those which implement the required domain.

The second part restricts the choice of implementation for the constraints by checking each value we will give for each parameter. The += operator is used to attach properties to a

```
@Dominion(Location = "/lib/src/", FileName = "problem1.hpp")
template ThisProblem() {
  provides IProblem pr;
  requires IPropVariable pvw, pvx, pvy, pvz;
  requires ISumCon scA, scB;

  check pvw->properties()  subsetof
      [(domainType, 'range'), (domainSize, 1)];
  check pvx->properties()  subsetof
      [(domainType, 'range'), (domainSize, 2)];
  check pvy->properties()  subsetof
      [(domainType, 'range')];
  check pvz->properties() subsetof
      [(domainSize, 3)];

  check scA->param(1) accepts (pvx  +=
    [(domainType, 'range'), (domainSize, 2), (length, 2)]);
  check scA->param(1) accepts
    (pvy  += [(domainType, 'range'), (length, 2)]);
  check scA->param(2) accepts (pvw  +=
    [(domainType, 'range'), (domainSize, 1)]);
  check scB->param(1) accepts pvz;
  check scB->param(2) accepts (pvw  +=
    [(domainType, 'range'), (domainSize, 1)]);
}
```

Listing 3: Specification of the Problem Component

component. We know in this particular problem that the domain of x will be $\{0, 1\}$. Regardless of the chosen implementation pvx for x, which may or may not make use of this restriction, any variable produced by pvx will have the domain $\{0, 1\}$.

C. Analyser

The work on the analyser tool is ongoing. This tool will generate a list of candidate solver architectures using the component library with the problem component as the seed and select the best one using artificial intelligence techniques.

Listing 4 shows a possible solver architecture produced by the analyser for the problem from Listing 1, including component instantiations and configurations (link statements). The template specifications are not shown due to lack of space.

```
system Solution {
  component vw  = ConstantVariableFactory();
  component vx  = BooleanVariableFactory();
  component vy  = GeneralVariableFactory();
  component vz  = GeneralVariableFactory();

  component conAf = AssignSumFactory();
  component conBf = AssignSumFactory();
  component problem = ThisProblem();

  link vw.var to problem.pvw;
  link vx.var to problem.pvx;
  link vy.var to problem.pvy;
  link vz.var to problem.pvz;

  link conAf.con to problem.scA;
  link conBf.con to problem.scB;
}
```

Listing 4: A Possible Solution Architecture

The reader is referred to existing literature on algorithm selection [8], [10] and algorithm portfolios [9] for more details on possible techniques for use by the analyser.

In practice, the variable and constraint components themselves have requires clauses, which leads to further components being instantiated. Given either a partial or complete solver, we can execute the check statements for each component to ensure the solver is valid.

D. Solver Generation

The architecture chosen by the analyser is given to the solver generator to create the target solver. This tool uses the location and file name information attached to each element in the architecture graph to find the component implementations required for the solver in the C++ component library. The main tasks of the solver are to:

- create a file for the main solver program,
- include the component files required by the chosen architecture
- instantiate the included components and parameterise them as appropriate, and
- generate code to set-up and begin the execution of the solver.

The solver generator performs a fairly straight-forward translation based on the decisions made earlier in the process.

VII. Conclusions and Future Work

We have provided an outline of a novel, architecture-driven approach to generating constraint solvers that are optimised for a given problem. The generative approach demands an expressive ADL, such as Grasp, able to capture different types of compatibility requirements among architectural elements.

In addition to completing the remaining tools and evaluating the generative framework against existing solvers, numerous interesting avenues remain for further work. These include extending the role of software architectures in assisting the tasks of analysis and execution monitoring.

Acknowledgments

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