# Scalable Normalization of Heap Manipulating Functions 

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heap

| $\mathrm{a0}$ | $\mathrm{v0}$ |
| :---: | :---: |
| a 1 | v 1 |
| a 2 | v 2 |
|  |  |
|  | $\circ$ |
|  |  |
|  |  |
|  |  |
| aK | v |
| aM | vK |
| aN | vN |



## Non-Interference


(defthm read-over-write-non-interference (implies
(not (equal a1 a2))
(equal (read a1 (write a2 v x)) (read a1 x))))

## Read/Write Towers



## Read/Write Towers

(defun write_3 (avr)
(defun read_3 (a r)
(write_2 a v r)) (read_2 a r))
(defun write_2 (avr)
(defun read_2 (a r)
(write_1 a v r))
(read_1 a r))
(defun write_1 (avr)
(write avr))
(defun read_1 (a r)
(read ar))
(write avr)
read-over-write-non-interference
(read ar)

## Read/Write Towers

(defun write_3 (avr)
(defun read_3 (a r)
(write_2 a v r)) (read_2 a r))
(defun write_2 (avr)
(write_1 a v r))
(defun read_2 (a r)
(read_1 a r))


## Read/Write Towers

(defun write_3 (avr)
(defun read_3 (ar)
(write_2 a v r)) (read_2 a r))
(defun write_2 (avr)
(write_1 a v r))
(defun read_2 (a r)
(read_1 a r))


## Read/Write Towers

(defun write_3 (avr)
(write_2 a v r))
(defun read_3 (ar) (read_2 a r))
(defun write_2 (avr)
(write_1 a v r))
(defun read_2 (ar)
(read_1 a r))


## Read/Write Towers

(defun write_3 (avr)
(defun read_3 (a r)
(write_2 a v r)) (read_2 a r))


## Read/Write Towers

(defun write_3 (avr)
(write_2 a v r))
(defun read_3 (ar) (read_2 a r))


## Read/Write Towers



## Non-interference in complex systems

- Complex Systems
- Hierarchical Design
- Build larger components from many simpler components
- Compositional Verification Methodology Essential
- Specify behavior once (locally)
- Use behavior many times (globally)
- Non-interference
- Not a complex property
- Number of theorems is quadratic in total number of components
- Standard Approach
- Articulate property between every component
- Not Compositional
- Doesn't scale


## And now for something completely different:

- Congruence-based Rewriting
- Built-In to ACL2
- Treats Certain Predicate Relations "just like equality"
- Use Relations to Define Rewrite Rules
- Provides Strong Normalization
- (Near) Minimal Representations
- Congruence-based Rewriting
- More powerful than rewrite rules
- More scalable than syntactic techniques (:meta / bind-free)
- Scalable
- Defined Locally
- Used Globally


## Rewriting Context

- Obviously (cons $x$ (cons $\times \mathrm{y}$ )) is not equal to (cons x y),

$$
\begin{aligned}
& (\text { cons } x(\text { cons } x y)) \\
& (\text { cons } x \text { y) }
\end{aligned}
$$

- But they are equivalent in "the second argument of member"

> (defthm member-cons-duplicates
> (iff (member a (cons x (cons x y)))
> $\quad($ member a (cons x y))))

- So we can replace one with the other in that context


## Defining a Rewriting Context

- ACL2 Generalizes this notion
- "the second argument of member"
- Uses Equivalence Relations
- Formalize essential properties of "the second argument of member"

```
(defun set-equiv (x y)
    (if (consp x)
        (and (member (car x) y)
            (set-equiv (cdr x) (remove (car x) y))
        (not (consp y))))
```

- Formally Introduced in ACL2 via defequiv
- (defequiv set-equiv)
- Associates equivalence relation with a rewriting context
- Rewrite rules employing equivalence relations

> (defthm set-equiv-cons-cons-driver
> (set-equiv (cons $\times($ cons $x$ y)) (cons $x$ y)))

- Does not rewrite set-equiv to true
- Replaces (cons $x$ (cons $x y$ )) with (cons $x y$ )
- In a set-equiv rewriting context
- Driver Rules
- Concise, Automatic, Unconstrained
- Enhanced Normalization


## Congruences

- Driver Rules
- Only Applied in specific rewriting contexts
- Congruence Rules
- Establish rewriting contexts
- Indicate when it is sound to use specified equivalence relations
(defthm set-equiv-implies-iff-in-2
(implies
(set-equiv $x$ y)
(iff (member a x) (member a y)))
:rule-classes (:congruence))


## Congruence-based Rewriting: Synopsys

- Rewriting contexts
- Characterized by equivalence relations
- Driver Rules
- Apply context-sensitive simplifications
- Congruence Rules
- Chain from one context to another
- Congruence-based Rewriting
- More powerful than rewrite rules
- More scalable than syntactic techniques
(defequiv set-equiv)
(defthm set-equiv-cons-cons-driver
(set-equiv (cons x (cons x y))
(cons xy)))
(defcong set-equiv iff (member a x) 2)
(defcong set-equiv set-equiv (cons a x) 2)


## Nary Congruences

- Nary Library
- Extends ACL2 congruence capabilities
- Enables parameterized equivalence relations and congruences
- Used to define parameterized rewrite rules



## Non-Interference as a Congruence

- Non-interference properties can be expressed via parameterized congruences
- Given an appropriate equivalence relation
- Inherits Congruence Properties
- Provides Strong Normalization
- (Near) Minimal Representations
- Scalable
- Defined Locally
- Used Globally

|  |  | $\stackrel{\text { use-equiv }}{\longleftrightarrow}$ |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  |  | aY | vY |
| al | vl |  | al | vl |
|  |  |  |  |  |
| aJ | vJ | (al, aJ, aK) | aJ | vJ |
|  |  | $\uparrow$ |  |  |
| aX | vX |  |  |  |
| aK | vK | $1$ | aK | vK |
|  |  |  |  |  |

(if (consp list)
(and (equal (read (car list) x) (read (car list) y))
(use-equiv x y (cdr list)))))

(defthm read-use-cong
(implies
(use-equiv x y (list a))
(equal (read a x)

> Nary congruence rule (read a y))))

(defthm write-use-elim (implies


X
uset)))

(defthm write-use-cong
(implies
(use-equiv $x$ y uset)
(use-equiv (write a $v x$ )

> Nary congruence rule
(write a $v$ y) uset)))

(defthm write-use-cong
(implies
(use-equiv x y (remove a uset)) $\longleftarrow$ (use-equiv (write a $v x$ )
(write a $\vee \mathrm{y}$ ) uset)))

## Local Characterization

(defthm read-use-cong (implies (use-equiv x y (list a)) (equal (read a x) (read a y))))
(defthm write-use-elim (implies (not (member a uset)) (use-equiv (write a $v x$ ) X uset)))
(defthm write-use-cong (implies
(use-equiv x y (remove a uset))
These three theorems characterize the non-interference properties of read and write operations via use-equiv

These three theorems are sufficient to characterize the non-interference properties of any function defined in terms of read and write.

Local characterization and global application: properties essential for scalable non-interference
(write a $\vee \mathrm{y}$ ) uset)))

## Example Application

(defthm read-use-cong (implies
(use-equiv x y (list a))
(equal (read a x)
(read a (write a v3 x)))))

## Example Application

(defthm read-over-write-normalization

| (list a) | (implies | (defthm write-use-elim |
| :---: | :---: | :---: |
|  | (not (member a (list b c d)) | (implies |
|  | (equal (read a (write bv1 | (not (member a uset)) |
|  | (write c v2 | (use-equiv (write av x ) |
|  | (write a v3 | $\underline{ }$ ( |
|  | (write d v4 (write $a v 5 x))$ )) | uset)) |

(read a (write a v3x)))))

## Example Application

(defthm read-over-write-normalization
(implies
(not (member a (list b c d))) (equal (read a
(write c v2 (write a v 3
(write d v4 (write a v5 x)))))
(read a (write a v3x)))))

## Example Application

(defthm read-over-write-normalization
(list a) (implies
(not (member a (list b c d)))
(equat (read a

$(\operatorname{read} a($ write av3 $x)))))$
(defthm write-use-cong (implies
(use-equiv x y (remove a uset))
(use-equiv (write a $v x$ )
(write a $\vee \mathrm{y}$ )
uset)))

## Example Application

(defthm read-over-write-normalization
(implies
(not (member a (list b c d))) (equal (read a
(defthm write-use-elim (implies (not (member a uset))

(read a (write a v3x)))))

## Example Application

(defthm read-over-write-normalization (implies
(not (member a (list b c d))) (equal (read a
(defthm write-use-elim (implies (not (member a uset))
(use-equiv (write a $v x$ )
uset)))
(read a (write a v3x)))))

## Example Application

(defthm read-over-write-normalization (implies
(not (member a (list b c d))) (equal (read a

(read a (write a v3x)))))

## Tower Example

(defun write_3 (avr) (write_2 a v r))
(implies
(defun write_2 (a v r) (equal (read_i a (write_j b v1 (write_1 a v r))
(defun write_1 (a v r) (write avr))
(write a $\vee$ r)
(defthm read_i-over-write_x-normalization
(not (member a (list b c d)))
(write_k c v2
$($ write_x a v3
$($ write_y d v4
$($ write_z a v5 x) )) )) )
$($ read i a $($ write x a v3 x) $))))$
$($ write_x a v3
$($ write_y d v4
$($ write_z a v5 x) )) )) )
$($ read i a $($ write x a v3 x) $))))$
(write_z a v5 x))))))
(read_i a (write_x a v3 x)))))
(defun read_3 (ar) (read_2 a r))
(defun read_2 (ar) (read_1 a r))
(defun read_1 (a r) (read ar))
(read ar)

## Read/Modify/Write

(defun move (rptr wptr r) (write wptr (read rptr r) r))
(defthm move-use-cong
(implies
(use-equiv $x$ y (cons rptr (remove wptr uset))
(use-equiv (move rptr wptr $x$ )
(move rptr wptr y) uset)))
(defthm move-use-elim (implies
(not (member wptr uset))
(use-equiv (move rptr wptr $x$ )
X
uset)))

## Crawlers



## Parameterized Congruence for Non-Interference

- Non-interference properties can be expressed via parameterized congruences
- use-equiv
- Inherits Congruence Properties
- Provides Strong Normalization
- (Near) Minimal Representations
- Scalable
- Defined Locally
- Used Globally

