Verifying NoC communication architectures with ACL2

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Platform-Based Design and Networks on a Chip

- Platform-Based Design:
  - Re-use of parameterized modules (*Intellectual Properties*)
  - High-level of abstraction
  - *Communication*-centric: from buses to networks
- System Verification:
  - Proof of each component
  - Proof of their interconnection
- State-of-the-Art:
  - Model checking or theorem proving of instances of systems
  - Often at RTL and below
- The GeNoC Approach
  - A generic model for reasoning about NoCs

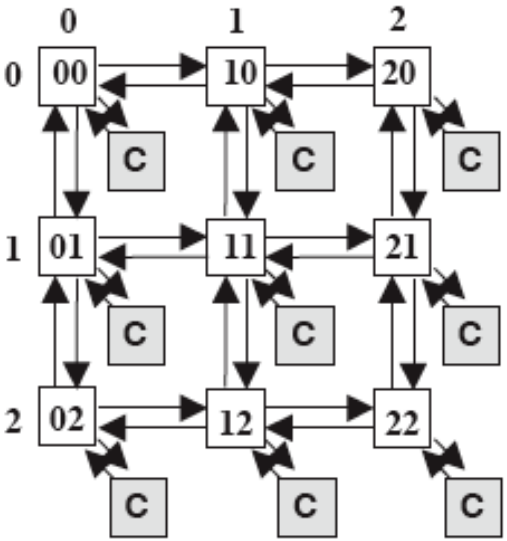
The GeNoC Model

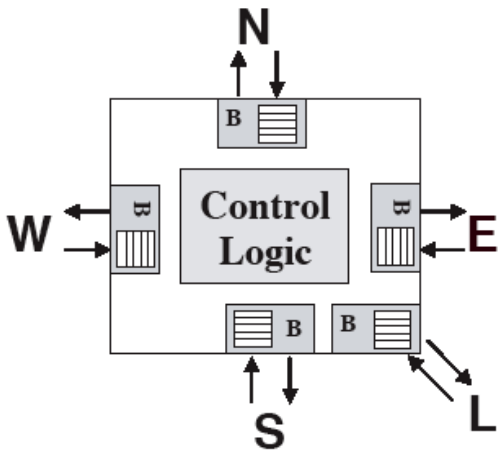
- Meta-model
  - Network topology and size
  - Routing algorithms
  - Switching techniques
- High-level of Abstraction
  - Abstract view of *Transport* (4) and *Network* (3) of the OSI model
- Encoded in the ACL2 theorem prover
  - Functional formalism
  - Parameterized proofs

An NoC Example: the Hermes NoC

- Result from cooperation between Catholic University of Rio Grande do Sul and LIRMM
- Public distribution
- FPGA implementation
- Protocol: credit based control flow
- Routing: deterministic minimal XY routing in a 2D-mesh
- Scheduling: wormhole switching

XY Routing and Wormhole Switching





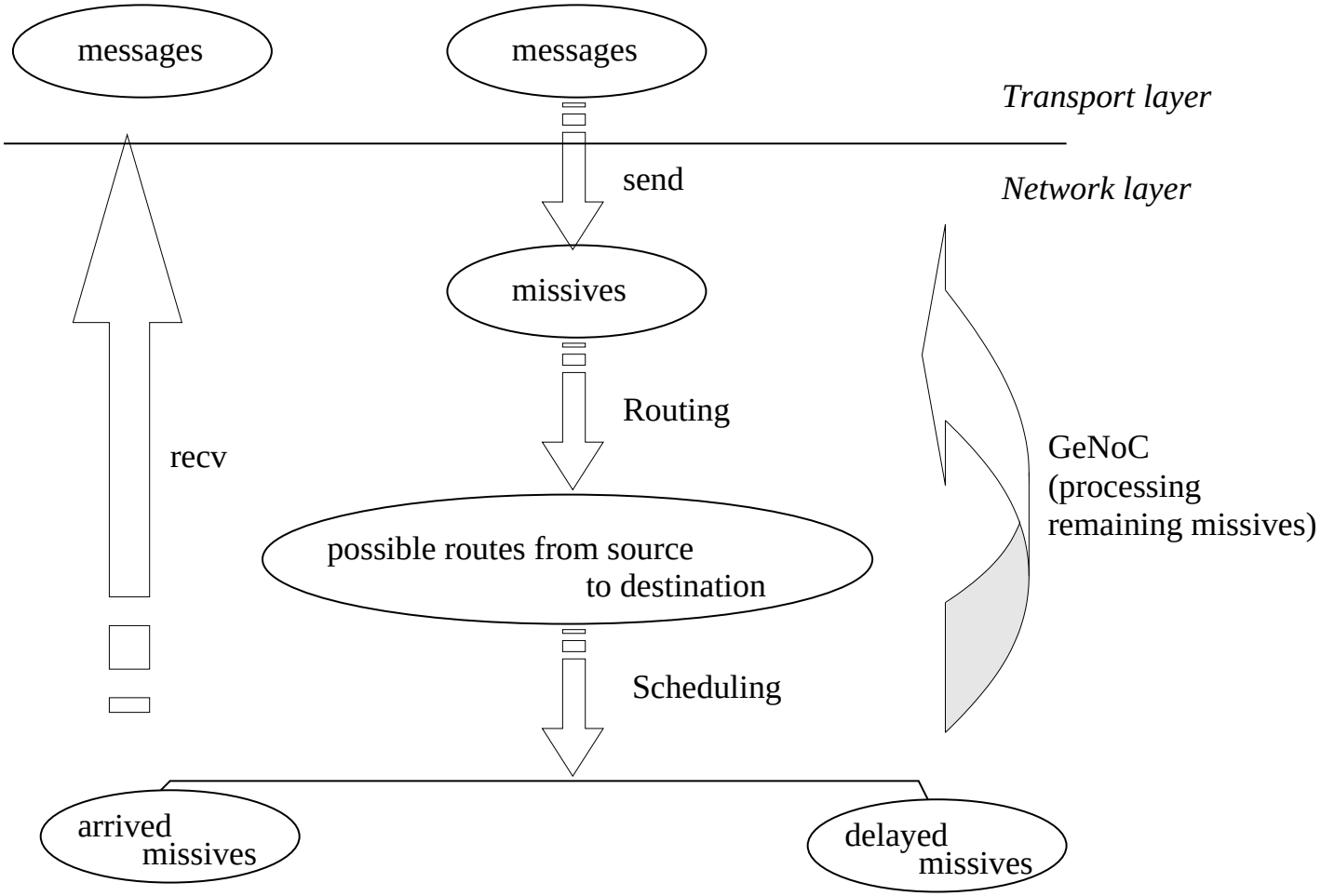
- XY minimal deterministic routing
- Frame structure:
  - Header flit (Route Information)
  - Control flit (Number of Flits)
  - Data flits (Payload)

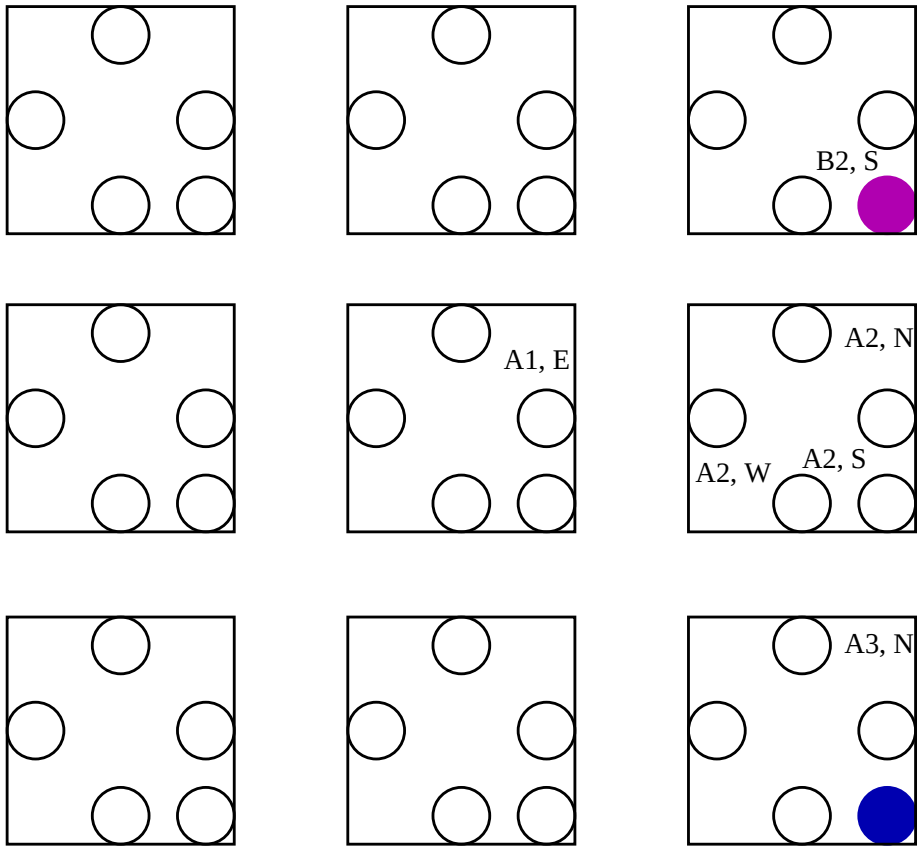
GeNoC and Hermes

- Initial model
  - Nodes associated with coordinates
  - Several messages cannot occupy a node simultaneously
  - Messages are atomic
  - No explicit notion of time
- Extensions [NOCs 2007]
  - Nodes have coordinates as well as ports
  - A node can be occupied by several messages
  - Support for non-atomic messages (flits)
- Application to Hermes
  - Modeling and validation using the extensions
  - Simulation in ACL2 $\neq$ Simulation VHDL
  - Messages blocked at origin if conflicts

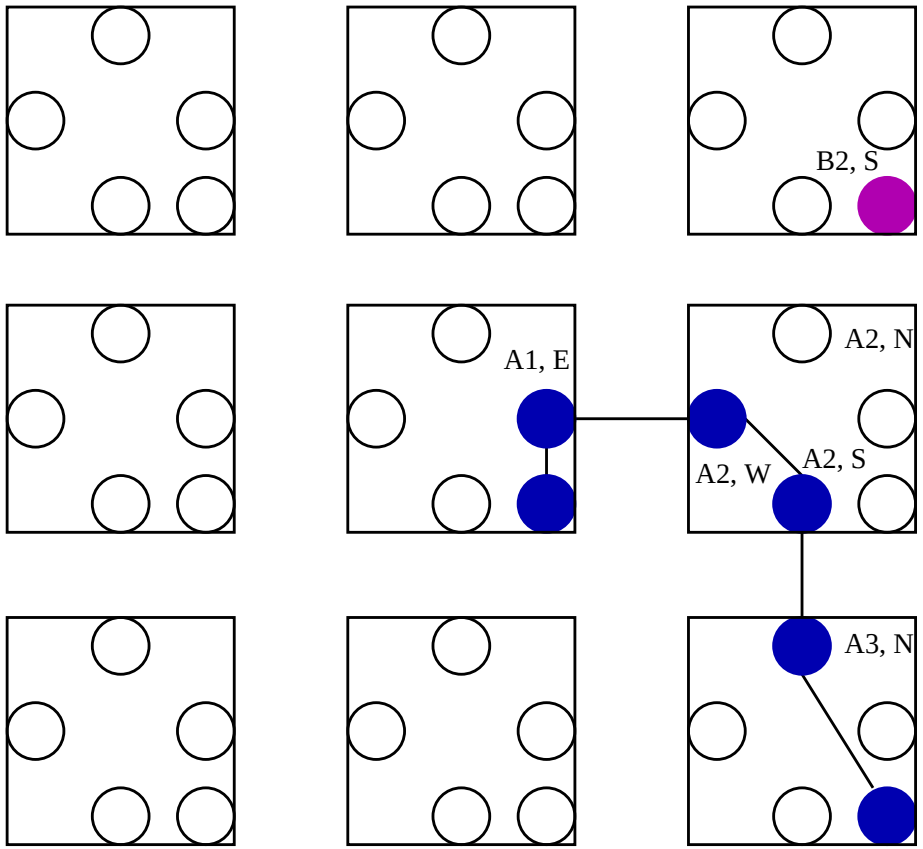
GeNoC Definition

```
( defun GeNoC_t (M NodeSet att TrLst)
  (declare (xargs :measure (SumOfAttempts att)))
  (if (zp (SumOfAttempts att))
      (mv TrLst M)
    (let ((V (Routing M NodeSet))) ;; comp. all routes
      (mv-let (Scheduled Delayed newAtt)
              ;; scheduled = reach dest.
              (Scheduling V att)
              (GeNoC_t (ToMissives Delayed) NodeSet newAtt
                       (append Scheduled TrLst))))))
```


messages
messages
Transport layer
send
Network layer
missives
Routing
recv
GeNoC
(processing
remaining missives)
possible routes from source
to destination
Scheduling
arrived
missives
delayed
missives



- Scheduled missives reach destination in one computation
- If conflicts, blocking missives stay home, they do not stay blocked in the network
- All missives sent at the same time



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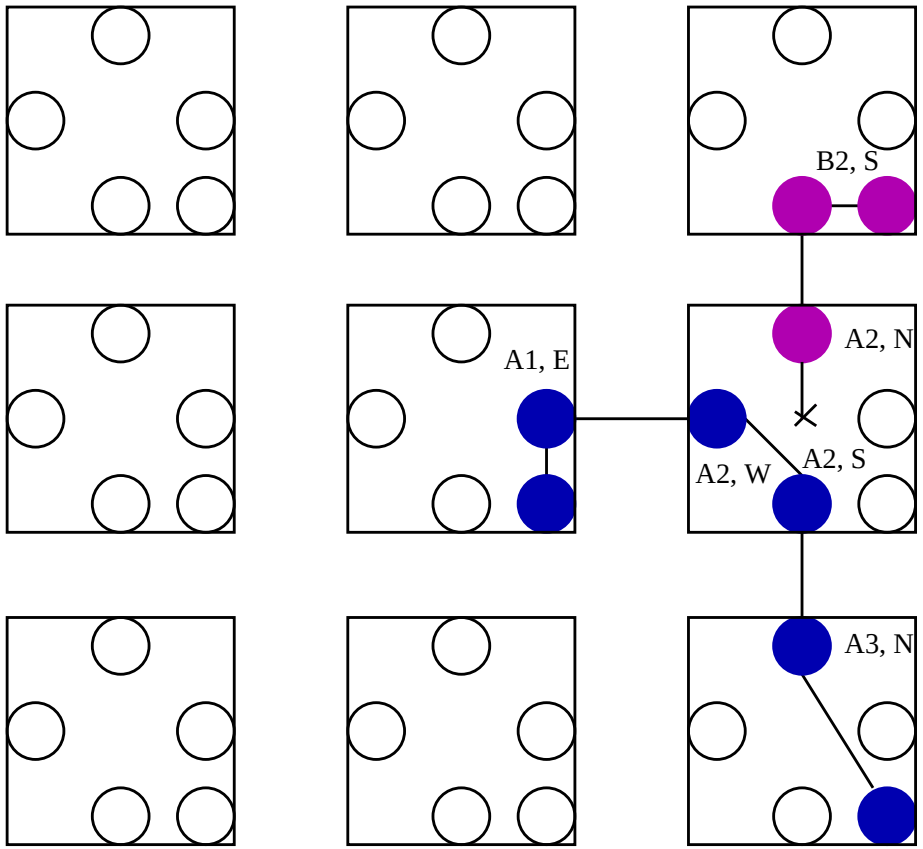

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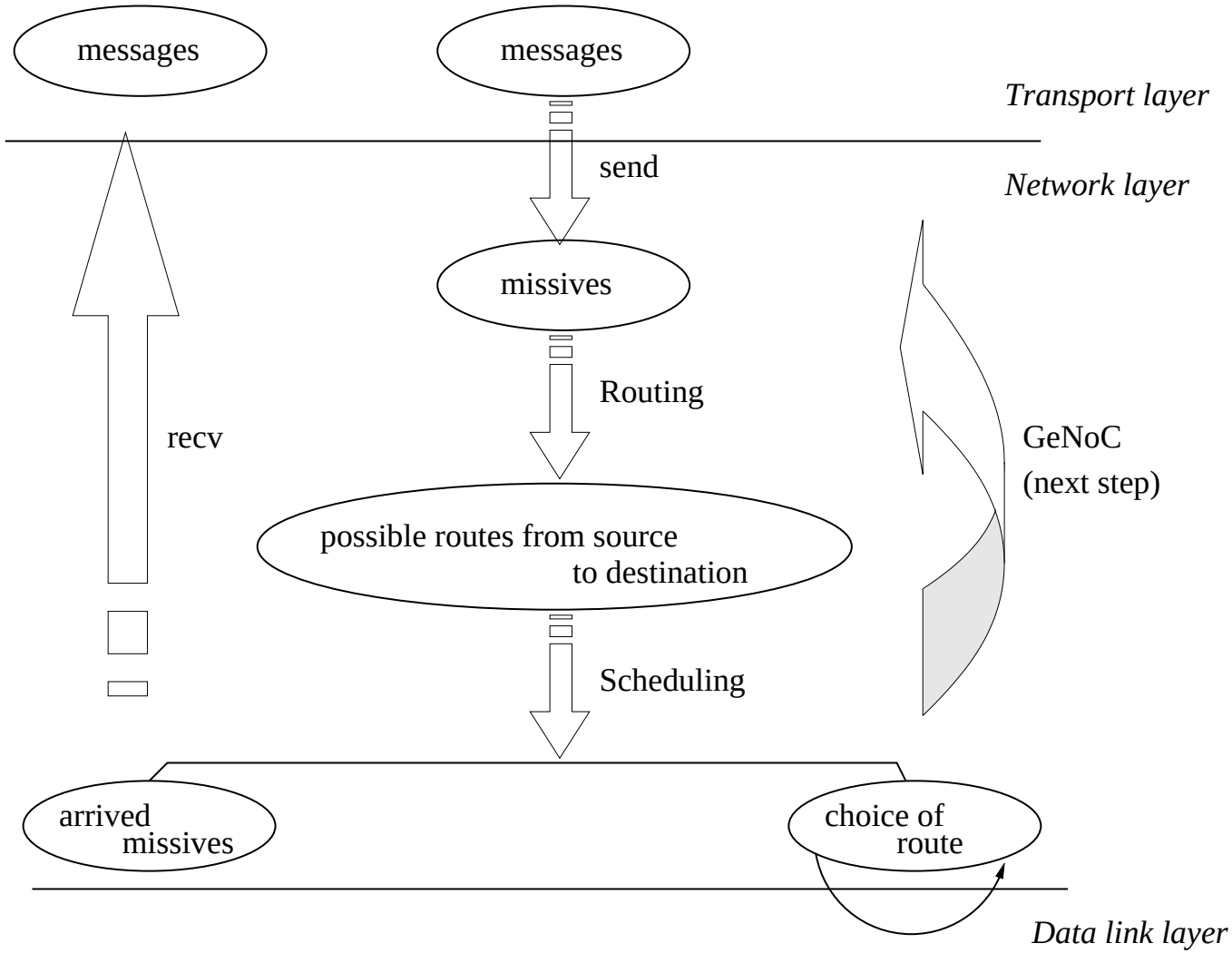
messages
messages
Transport layer
send
Network layer
missives
Routing
recv
GeNoC
(next step)
possible routes from source
to destination
Scheduling
arrived
missives
choice of
route
Data link layer

GeNoC: The Future

- Simulation ACL2 = Simulation VHDL
  - Improved step-by-step simulation
  - Global network state
- Deadlock avoidance
  - Protocol and structural
  - Control flow mechanism (credits)
- Link with RTL !
  - Data link layer (2) of OSI model