### A Generalized Solution for the While Challenge

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## Bill Young's "While Challenge"

### Introduce the following equation in ACL2

Yesterday's talk: Challenge answered by John Cowles and Dave Greve

> Additionally requires that run is strict in (btm)

### Kaufmann's "Generalized While" Challenge

Implement a macro for defining operational semantics of languages with unbounded while loops

# Show that a more general reflexive equation can be introduced with ACL2

This talk reports progress in answering Kaufmann's challenges.

### **Our Results**

Implement a macro for defining operational semantics of languages with unbounded while loops

- Developed a macro definiterpreter to introduce such semantics
- Show that a more general reflexive equation can be introduced with ACL2
  - Introduced the suggested equation about run given encapsulated functions test1, test2, finish, btm, etc.

### **Basic Approach**

#### First define a "clocked version" of run.

#### Then eliminate clk using quantification.

#### Essentially a formalization of Cowles' proof in an abstract setting.

### Macro for Language Interpreter

- Young's equation can be introduced by appropriate functional instantiation of test1, test2, dst1, stp, etc.
  - Cowles [private communication] showed the functional instance necessary.
- □ My macro definterpreter automates the functional instantiation and can introduce languages with unbounded while loops.

□ Provides some executability support via mbe construct.

### A Sneak Peek at Macro

```
(definterpreter run stmt mem
    :op-field (op stmt)
    :bottom nil
    :executable t
    :verify-quards nil
    :vanilla-interpreter (((:name skip)
                            (:interpreter mem))
                           ((:name assign)
                            (:interpreter (run-assignment stmt
mem))))
    :sequence ((:name sequence)
               (:argl (argl stmt))
               (:arq2 (arq2 stmt)))
    :conditional ((:name if)
                   (:test (zerop (eval-expr (arg1 stmt)
mem)))
                   (:tbr (arg3 stmt))
                   (:fbr (arg2 stmt)))
    :while ((:name while)
            (:test (zerop (eval-expr (arg1 stmt) mem)))
            (:body (arg2 stmt))))))
```

# Coming Up

Cowles showed that a number of reflexive equations can be introduced by functional instantiation of the generic theorem.

- Stay tuned for the next talk
- □ We are developing a macro defreflexive to automate introduction of such equations.
  - □A very preliminary implementation is available.

### **Details and Request**

For details please see **books/workshops/2007/cowles-et-al/support/ray/** 

I will appreciate any comments, and in particular interface suggestions for the macros.