Experience implementing a performant category-theory library in Coq

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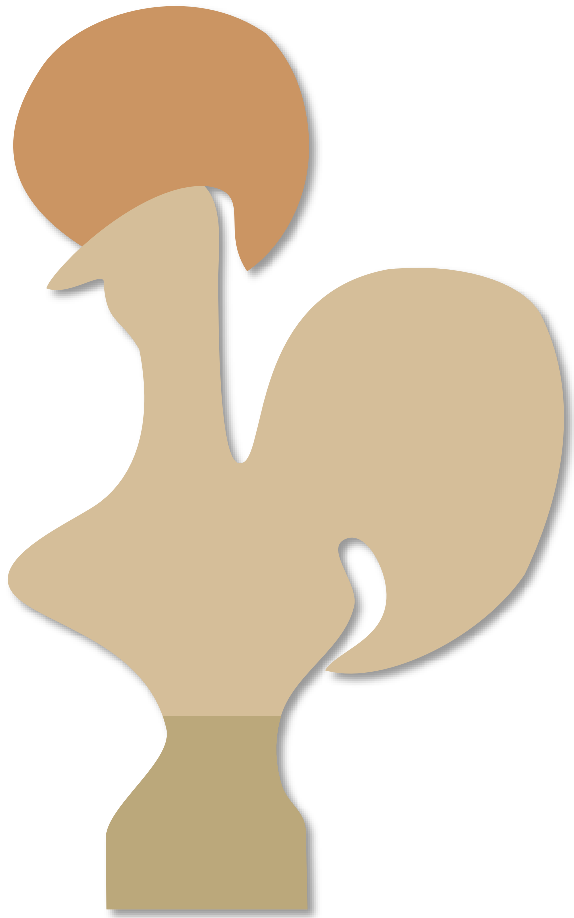
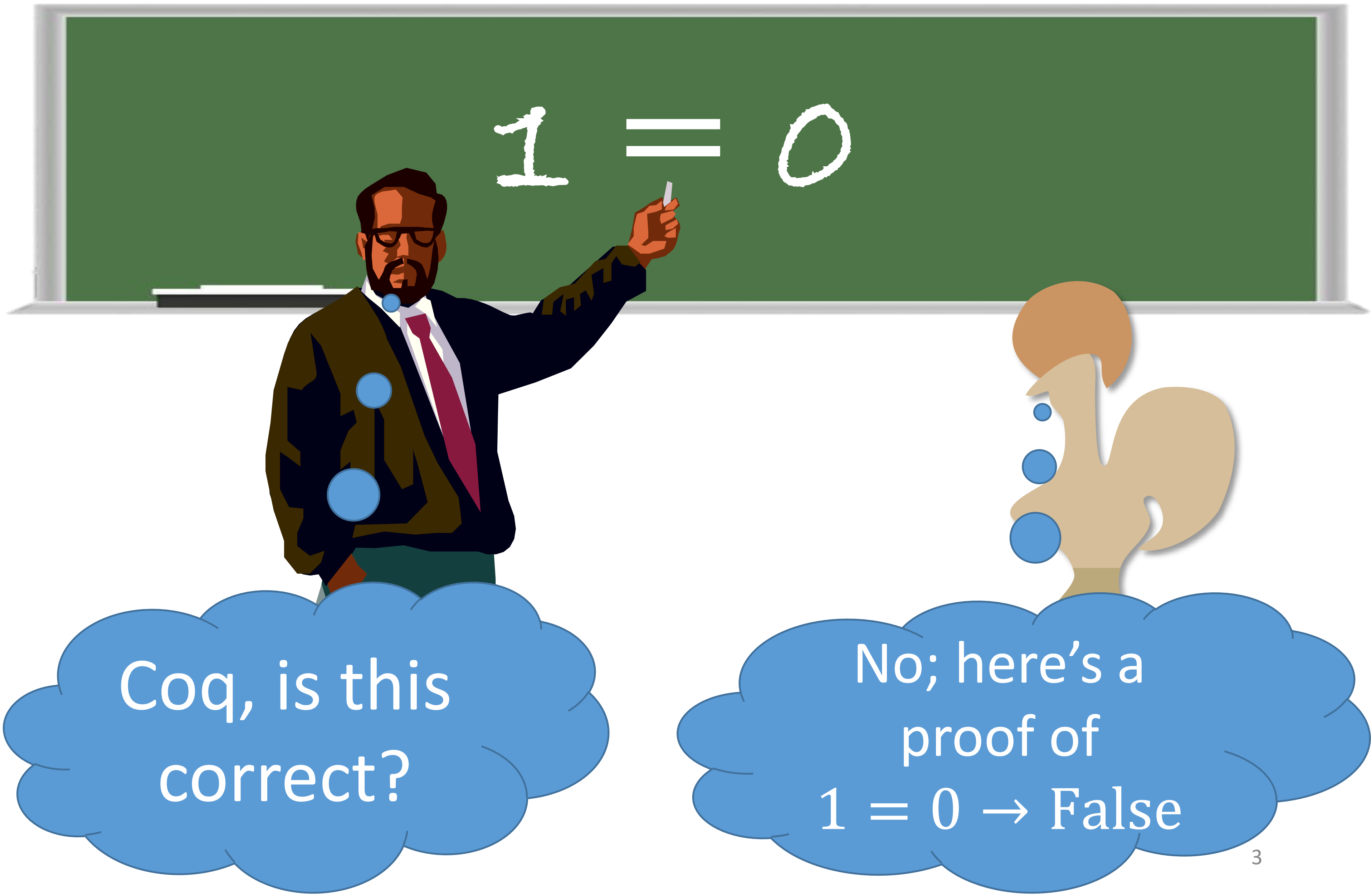
Massachusetts Institute of Technology

How should theorem provers work?



How theorem provers should work:



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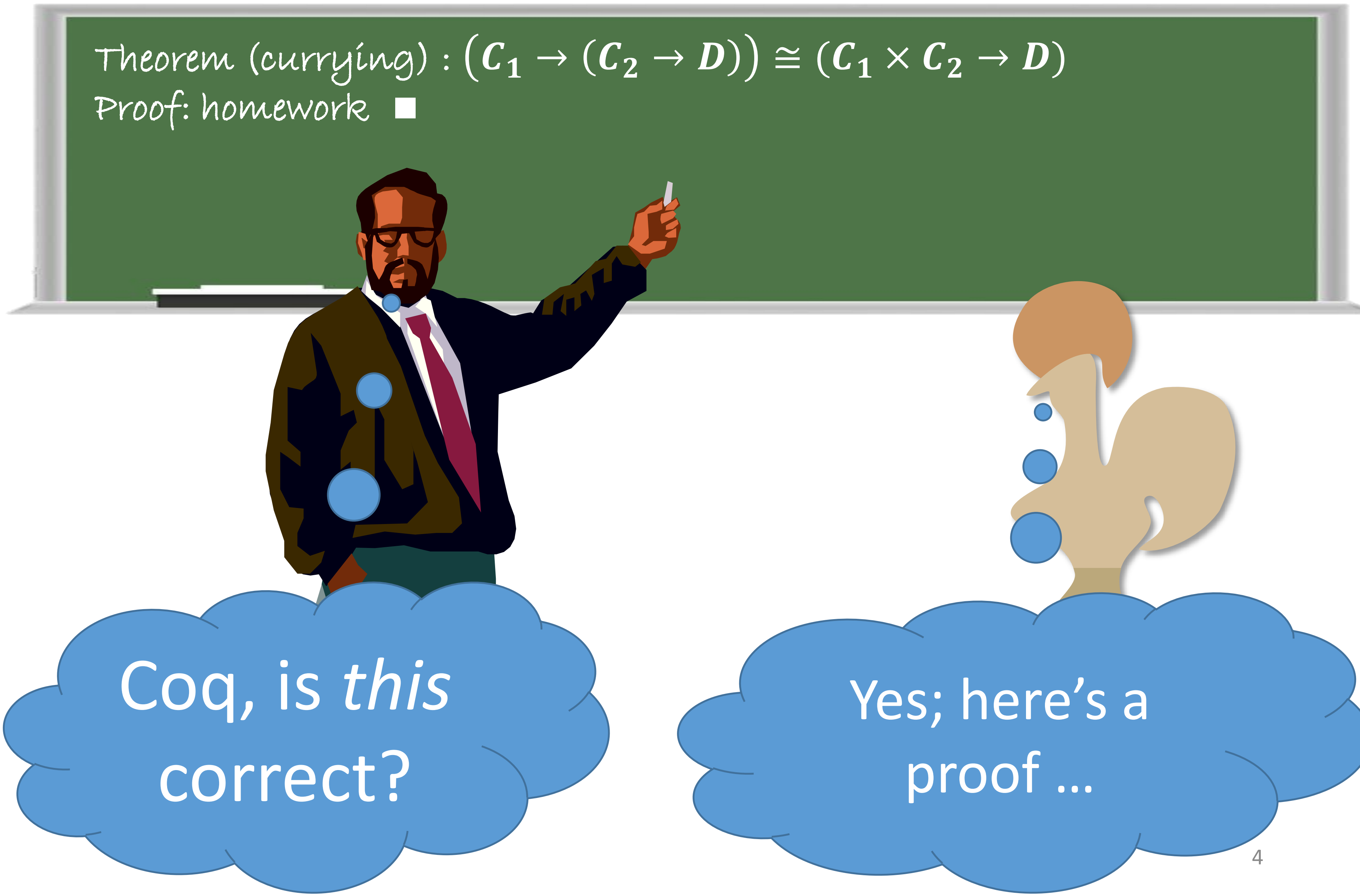

How theorem provers should work:

Theorem (currying) : $\left(\boldsymbol{C_1} \rightarrow (\boldsymbol{C_2} \rightarrow \boldsymbol{D})\right) \cong (\boldsymbol{C_1} \times \boldsymbol{C_2} \rightarrow \boldsymbol{D})$

Proof: homework ■

Theorem currying : $\left(C_1 \ \rightarrow (C_2 \rightarrow D)\right) \cong (C_1 \times C_2 \rightarrow D)$.

Proof.

  trivial.

Qed.

How theorem provers should work:

Theorem (currying) : $\big(C_1 \to (C_2 \to D)\big) \cong (C_1 \times C_2 \to D)$

Proof: $\to$: $F \mapsto \lambda\,(c_1, c_2).\ F(c_1)(c_2)$; morphisms similarly

$\leftarrow$: $F \mapsto \lambda\, c_1. \lambda\, c_2. F(c_1, c_2)$; morphisms similarly

Functoriality, naturality, and congruence: straightforward. ■

Theorem currying : $\big(C_1 \ \to (C_2 \to D)\big) \cong (C_1 \times C_2 \to D)$.

Proof.

esplit.

{ by refine $(\lambda_\mathrm{F}\ (F \mapsto (\lambda_\mathrm{F}\ (c \mapsto F_\mathrm{o}\ c_1\ c_2))))$. }

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all: trivial.

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$(F\ G\ T \mapsto (\lambda_{\mathrm{T}}\ (c_1 \mapsto (\lambda_{\mathrm{T}}\ (c_2 \mapsto T\ (c_1, c_2))))))$. }

all: trivial.

Qed.

How theorem provers do work:

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≈ 0 s

Functoriality, naturality, and congruence: straightforward. ■

17 s

2m 46 s !!! (5 s, if we use UIP)

Theorem currying : $\left(C_1 \to (C_2 \to D)\right) \cong (C_1 \times C_2 \to D)$.

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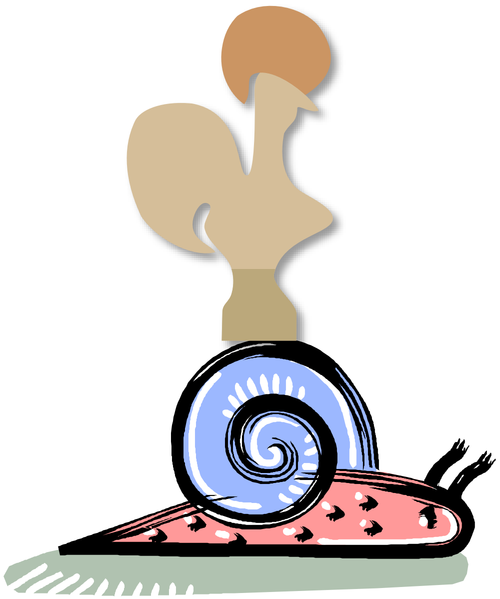
all: trivial.

Qed.

Performance is important!

If we're not careful, obvious or trivial things can be very, very slow.



Why you should listen to me

Theorem : You should listen to me.

Proof.

by experience.

Qed.

Why you should listen to me

Category theory in Coq: https://github.com/HoTT/HoTT (subdirectory theories/categories):

Concepts Formalized:

- 1-precategories (in the sense of the HoTT Book)
- univalent/saturated categories (or just categories, in the HoTT Book)
- functor precategories $C \to D$
- dual functor isomorphisms Cat → Cat; and $(C \to D)^{\mathrm{op}} \to (C^{\mathrm{op}} \to D^{\mathrm{op}})$
- the category Prop of (U-small) hProps
- the category Set of (U-small) hSets
- the category Cat of (U-small) strict (pre)categories (strict in the sense of the objects being hSets)
- pseudofunctors
- profunctors
  - identity profunction (the hom functor $C^{\mathrm{op}} \times C \to \mathrm{Set}$)
- adjoints
  - equivalences between a number of definitions:
    - unit-counit + zig-zag definition
    - unit + UMP definition
    - counit + UMP definition
    - universal morphism definition
    - hom-set definition (porting from old version in progress)
  - composition, identity, dual
  - pointwise adjunctions in the library, $G^E \dashv F^C$ and $E^F \dashv C^G$ from an adjunction $F \dashv G$ for functors $F: C \leftrightarrows D: G$ and E a precategory (still too slow to be merged into the library proper; code here)
- Yoneda lemma
- Exponential laws
  - $C^0 \cong 1$; $0^C \cong 0$ given an object in C
  - $C^1 \cong C$; $1^C \cong 1$
  - $C^{A+B} \cong C^A \times C^B$
  - $(A \times B)^C \cong A^C \times B^C$
  - $(A^B)^C \cong A^{B \times C}$
- Product laws
  - $C \times D \cong D \times C$
  - $C \times 0 \cong 0 \times C \cong 0$
  - $C \times 1 \cong 1 \times C \cong C$
- Grothendieck construction (oplax colimit) of a pseudofunctor to Cat
- Category of sections (gives rise to oplax limit of a pseudofunctor to Cat when applied to Grothendieck construction
- functor composition is functorial (there's a functor $\Delta: (C \to D) \to (D \to$

Presentation is **not** mainly about:

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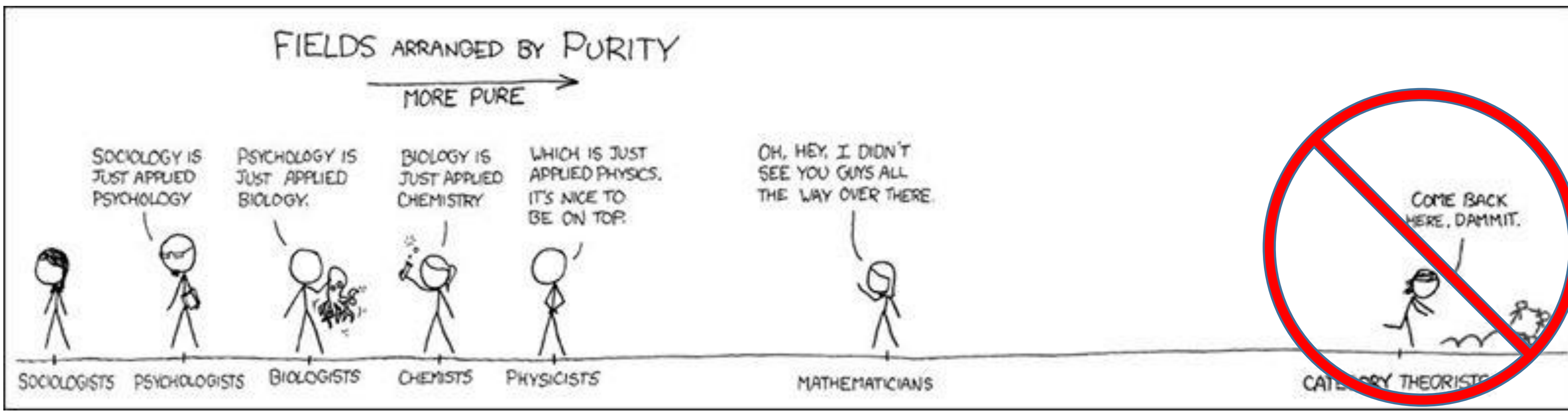
- category theory or diagram chasing



Cartoon from xkcd, adapted by Alan Huang

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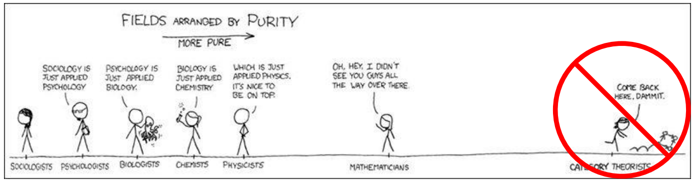

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- my library

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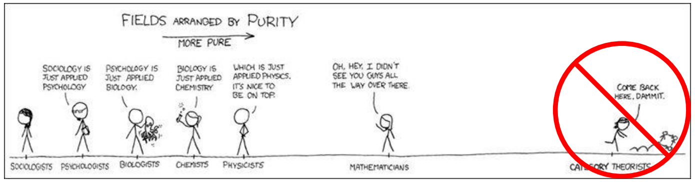

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- my library

- Coq

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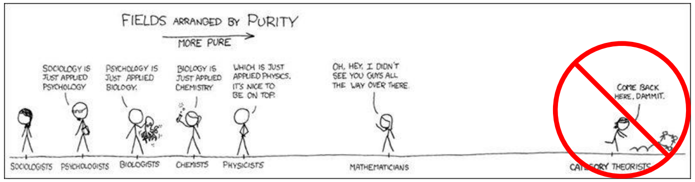
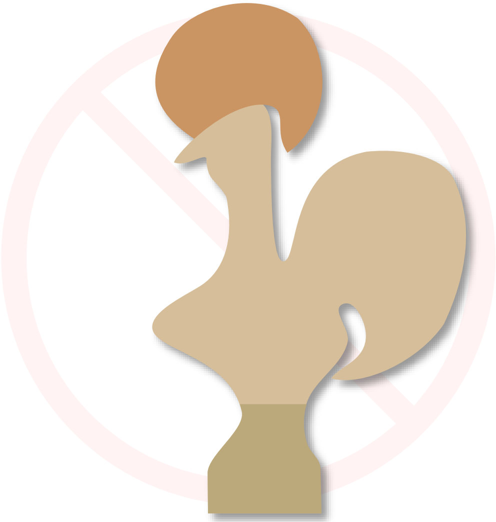


Cartoon from xkcd, adapted by Alan Huang

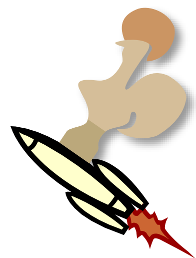
- category theory or diagram chasing
- my library

- Coq (though what I say might not always generalize nicely)



Presentation **is** about:

- performance



- the design of proof assistants and type theories to assist with performance

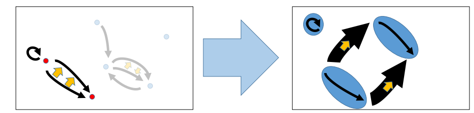
- the kind of performance issues I encountered

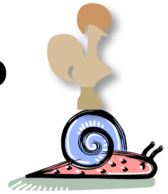
Presentation **is** for:

- Users of proof assistants (and Coq in particular)
  - Who want to make their code faster
- Designers of (type-theoretic) proof assistants
  - Who want to know where to focus their optimization efforts

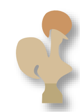
Outline

- Why should we care about performance?
- What makes theorem provers (mainly Coq) slow?
  - Examples of particular slowness

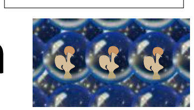




- For users (workarounds)
  - Arguments vs. fields and packed records
  - Proof by duality as proof by unification
  - Abstraction barriers
  - Proof by reflection



- For developers (features)
  - Primitive projections
  - Higher inductive types
  - Universe Polymorphism
  - More judgmental rules
  - Hashconsing



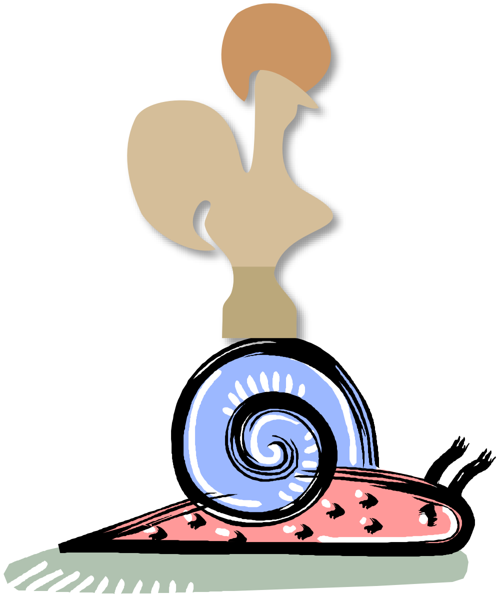
Universes image from Abell NGC2218 hst big, NASA, http://en.wikipedia.org/wiki/Abell_2218#mediaviewer/File:Abell_NGC2218_hst_big.jpg, released in Public Domain; Bubble from http://pixabay.com/en/blue-bubble-shiny-157652/, released in Public Domain CC0, combined in Photoshop by Jason Gross

Dam image from http://www.flickr.com/photos/gammaman/7803829282/ by Eli Christman, CC by 2.0

Fence image from http://www.picgifs.com/clip-art/playing-children/clip-art-playing-children-362018-689955/

Performance

- **Question:** What makes programs, particularly theorem provers or proof scripts, slow?



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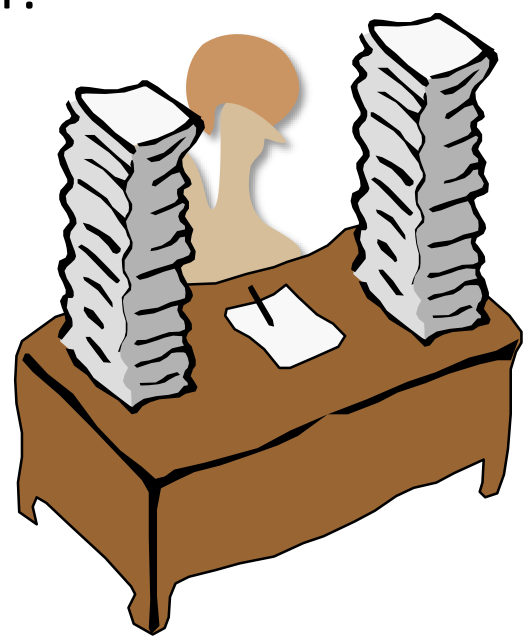

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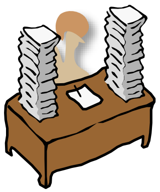

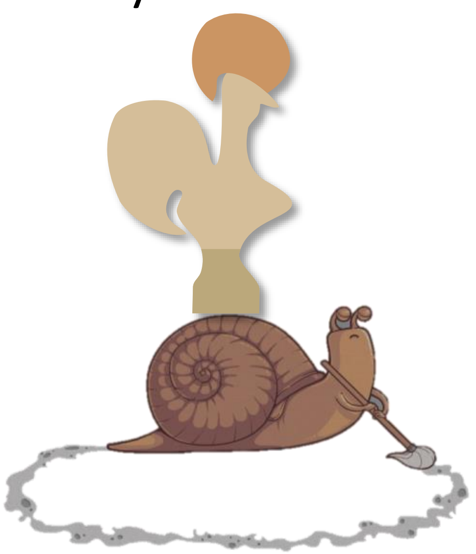

Snail from http://naolito.deviantart.com/art/Repetitive-task-258126598

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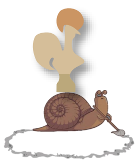

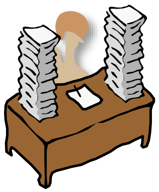

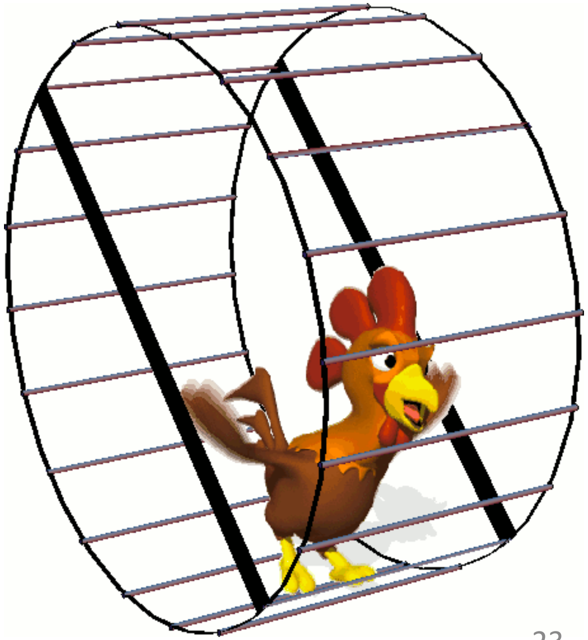

Running rooster from http://d.wapday.com:8080/animation/ccontennt/15545-f/mr_rooster_running.gif

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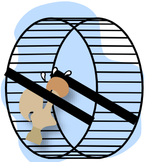
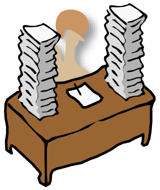
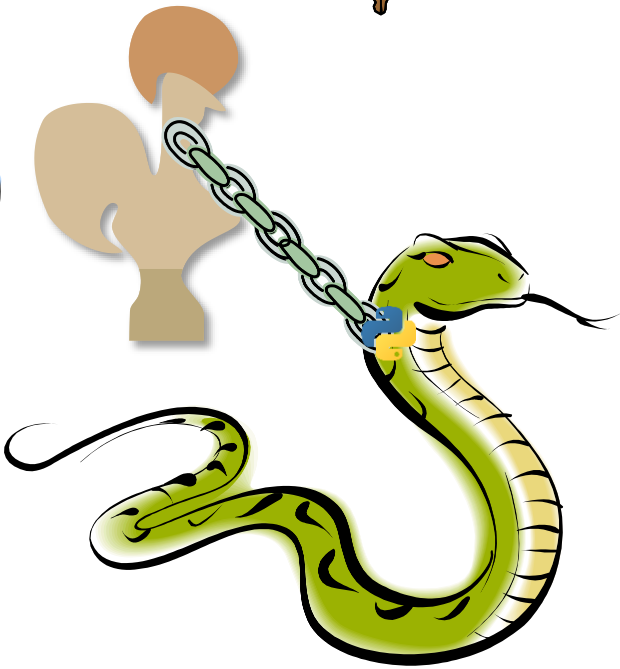
  
  - using a slow language when you could be using a quicker one





Proof assistant performance

- What kinds of things does Coq do?
  - Type checking
  - Term building
  - Unification
  - Normalization

Proof assistant performance (pain)

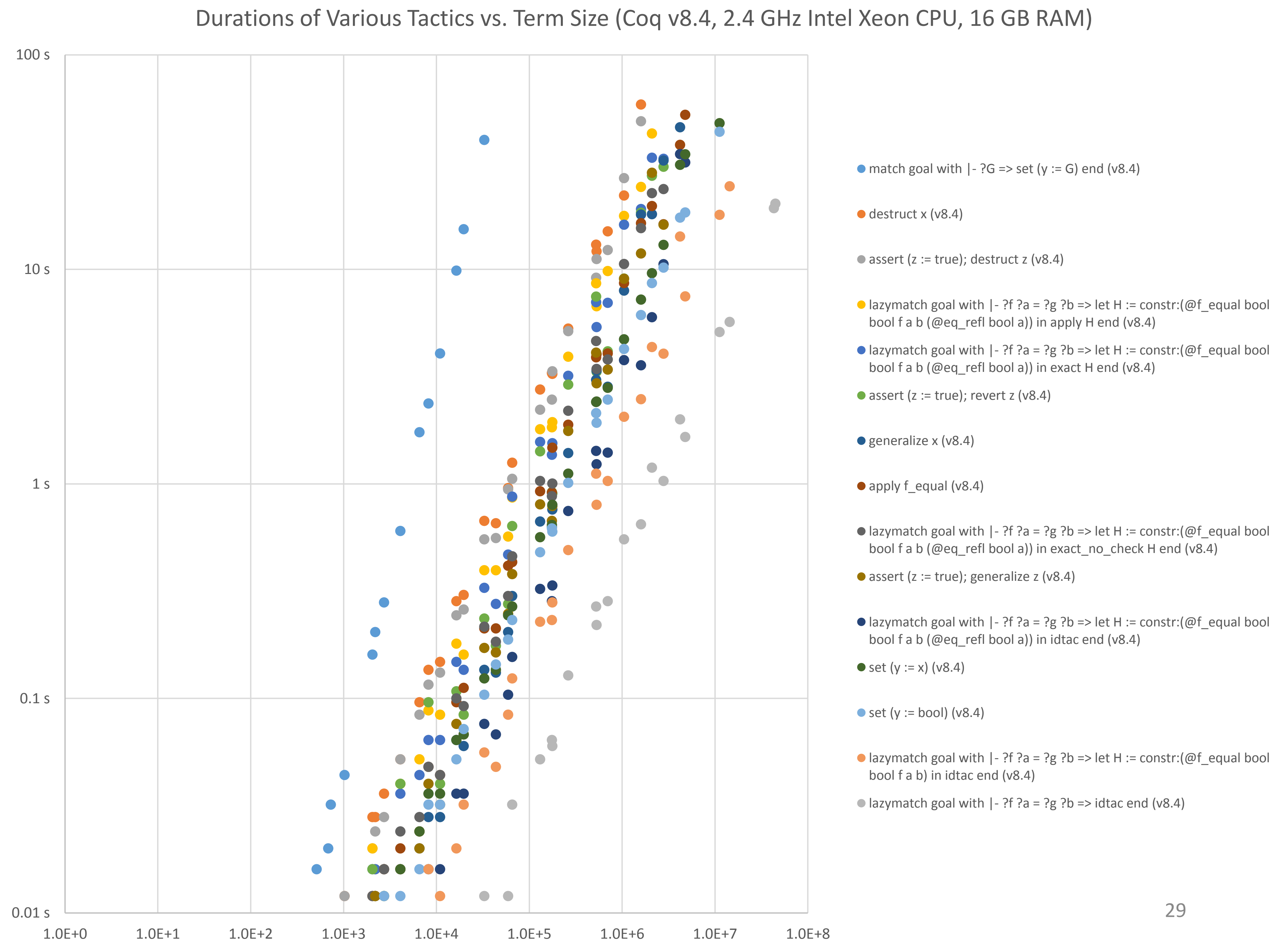
- When are these slow?
  - when you duplicate work
  - when you do work on a part of a term you end up not caring about
  - when you do them too many times
  - when your term is large

Proof assistant performance (size)

- How large is slow?

Proof assistant performance (size)

- How large is slow?
  - Around 150,000—500,000 words



Durations of Various Tactics vs. Term Size (Coq v8.4, 2.4 GHz Intel Xeon CPU, 16 GB RAM)
100 s
10 s
1 s
0.1 s
0.01 s
1.0E+0
1.0E+1
1.0E+2
1.0E+3
1.0E+4
1.0E+5
1.0E+6
1.0E+7
1.0E+8
match goal with |- ?G => set (y := G) end (v8.4)
destruct x (v8.4)
assert (z := true); destruct z (v8.4)
lazymatch goal with |- ?f ?a = ?g ?b => let H := constr:(@f_equal bool bool f a b (@eq_refl bool a)) in apply H end (v8.4)
lazymatch goal with |- ?f ?a = ?g ?b => let H := constr:(@f_equal bool bool f a b (@eq_refl bool a)) in exact H end (v8.4)
assert (z := true); revert z (v8.4)
generalize x (v8.4)
apply f_equal (v8.4)
lazymatch goal with |- ?f ?a = ?g ?b => let H := constr:(@f_equal bool bool f a b (@eq_refl bool a)) in exact_no_check H end (v8.4)
assert (z := true); generalize z (v8.4)
lazymatch goal with |- ?f ?a = ?g ?b => let H := constr:(@f_equal bool bool f a b (@eq_refl bool a)) in idtac end (v8.4)
set (y := x) (v8.4)
set (y := bool) (v8.4)
lazymatch goal with |- ?f ?a = ?g ?b => let H := constr:(@f_equal bool bool f a b) in idtac end (v8.4)
lazymatch goal with |- ?f ?a = ?g ?b => idtac end (v8.4)

Proof assistant performance (size)

- How large is slow?
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Do terms actually get this large?

Proof assistant performance (size)

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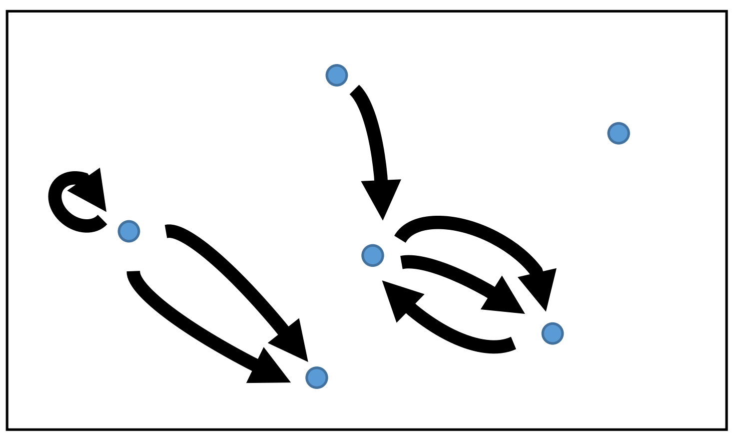
Do terms actually get this large?

YES!

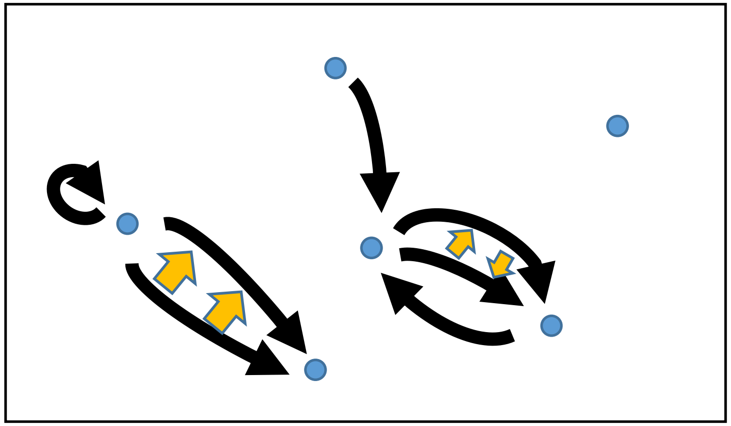
Proof assistant performance (size)

- A **directed graph** has:
  - a type of vertices (points)
  - for every ordered pair of vertices, a type of arrows



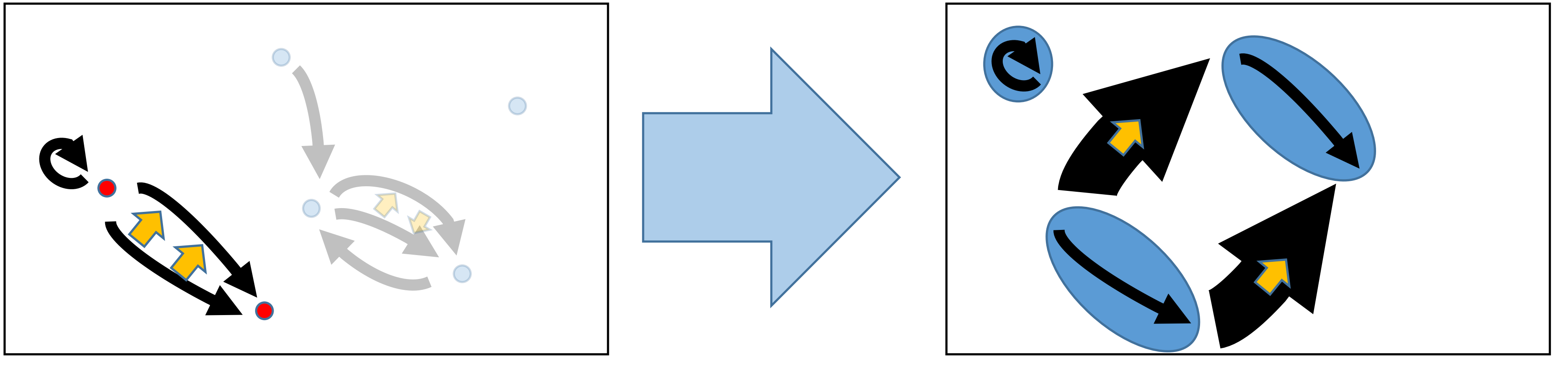
Proof assistant performance (size)

- A **directed 2-graph** has:
  - a type of vertices (0-arrows)
  - for every ordered pair of vertices, a type of arrows (1-arrows)
  - for every ordered pair of 1-arrows between the same vertices, a type of 2-arrows



Proof assistant performance (size)

- A **directed arrow-graph** comes from turning arrows into vertices:



Proof assistant performance (pain)

- When are these slow?
  - When your term is large
- Smallish example (29 000 words): Without Proofs:

$$\{|\ \ \mathrm{LCCM}_{\mathrm{F}} \coloneqq \_\backslash\_\mathrm{induced}_{\mathrm{F}}\ (m_{22} \circ m_{12});$$
$$\mathrm{LCCM}_{\mathrm{T}} \coloneqq \lambda_T\ (\lambda\ (c : d_2' \ / \ F) \Rightarrow m_{21}\ c.\beta \circ m_{11}\ c.\beta)\ |\} =$$
$$\{|\ \ \mathrm{LCCM}_{\mathrm{F}} \coloneqq \_\backslash\_\mathrm{induced}_{\mathrm{F}}\ m_{12} \circ \_\backslash\_\mathrm{induced}_{\mathrm{F}}\ m_{22};$$
$$\mathrm{LCCM}_{\mathrm{T}} \coloneqq \lambda_T\ (\lambda\ (c : d_2' \ / \ F) \Rightarrow m_{21}\ c.\beta \circ (d_1)_1\ \mathbb{I} \circ m_{11}\ c.\beta \circ \mathbb{I})\ |\}$$

Proof assistant performance (pain)

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$$(\Pi\mathrm{-pf}\ s_2\ (\lambda_T\ (\lambda\ (c : C) \Rightarrow m_{21}\ c \circ m_{11}\ c)$$
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$$(\circ_0\ \mathrm{-pf}\ (\lambda_T\ (\lambda\ (c : d_2\ /\ F) \Rightarrow \ldots$$
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$$
\begin{aligned}
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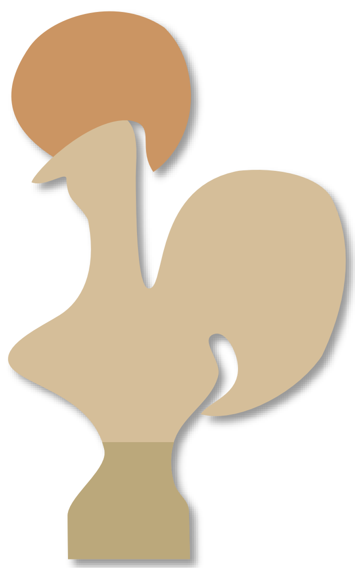
&\qquad\qquad (\circ_1 -\mathrm{pf}\ \ (\lambda_T\ (\lambda\ (c : d_2'\ /\ F) \Rightarrow m_{21}\ c.\beta)\ (\Pi{-}\mathrm{pf}\ d_2\ m_{21}\ m_{22})))\\
&\qquad\qquad\qquad (\lambda_T\ (\lambda\ (c : d_2'\ /\ F) \Rightarrow (d_1)_1\ \mathbb{I} \circ m_{11}\ c.\beta \circ \mathbb{I})\\
&\qquad\qquad\qquad\quad (\circ_1 -\mathrm{pf}\ (\lambda_T\ (\lambda\ (c : d_2'\ /\ F) \Rightarrow (d_1)_1\ \mathbb{I} \circ m_{11}\ c.\beta)\\
&\qquad\qquad\qquad\qquad (\circ_0 -\mathrm{pf}\ \ (\lambda_T\ \ (\lambda\ (c : d_2\ /\ F) \Rightarrow m_{11}\ c.\beta)\\
&\qquad\qquad\qquad\qquad\quad (\Pi{-}\mathrm{pf}\ s_2\ m_{11}\ m_{12}))\ \mathbb{I}))\ \mathbb{I})))\ |\}
\end{aligned}
$$

Proof assistant performance (fixes)

- How do we work around this?

Proof assistant performance (fixes)

- How do we work around this?
- By hiding from the proof checker!



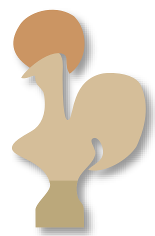
Fence from http://imgarcade.com/1/hiding-clipart/

Proof assistant performance (fixes)

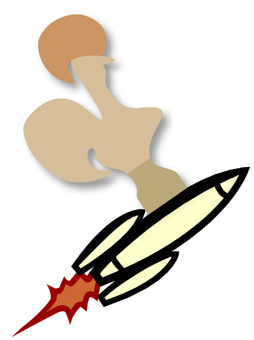
- How do we work around this?
- By hiding from the proof checker!
- How do we hide?

Proof assistant performance (fixes)

- How do we work around this?
- By hiding from the proof checker!
- How do we hide?
  - Good engineering



  - Better proof assistants

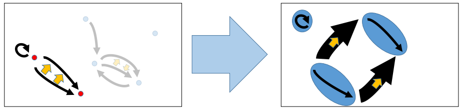


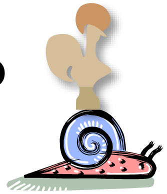
Proof assistant performance (fixes)

Careful Engineering

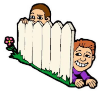
Outline

- Why should we care about performance?
- What makes theorem provers (mainly Coq) slow?
  - Examples of particular slowness





- **For users (workarounds)**
  - **Arguments vs. fields and packed records**
  - **Proof by duality as proof by unification**
  - **Abstraction barriers**
  - **Proof by reflection**



- For developers (features)
  - Primitive projections
  - Higher inductive types
  - Universe Polymorphism
  - More judgmental rules
  - Hashconsing

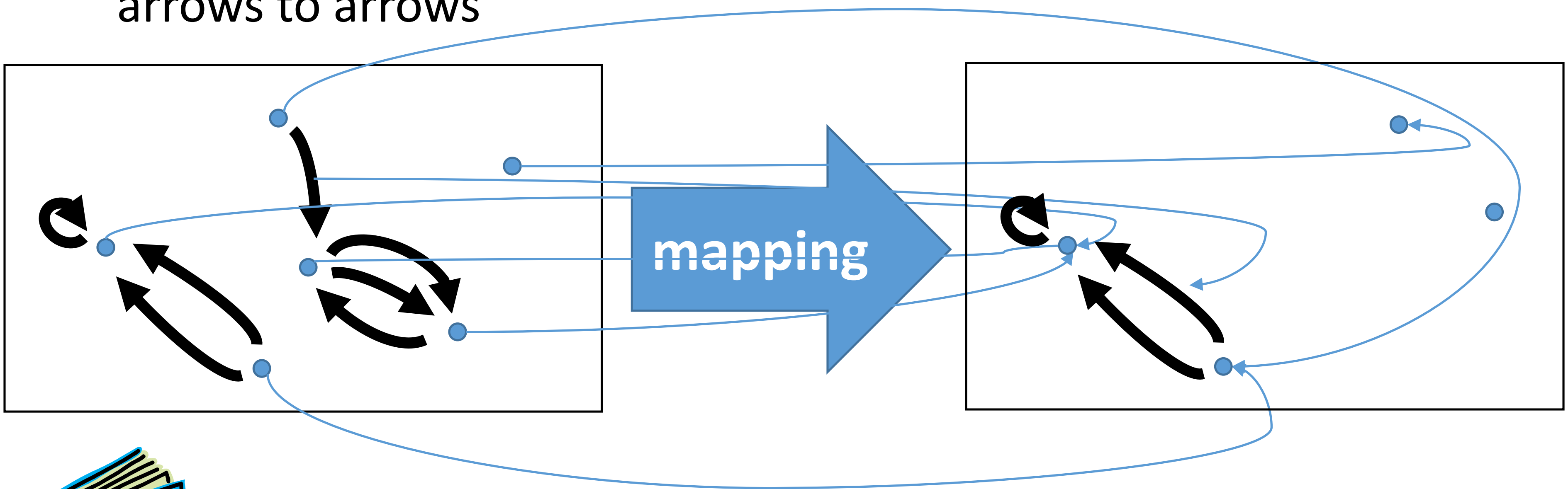
Proof assistant performance (fixes)

- How?
  - Pack your records!

Proof assistant performance (fixes)

- How?
  - Pack your records!

A **mapping of graphs** is a mapping of vetices to vertices and arrows to arrows



Proof assistant performance (fixes)

- How?
  - Pack your records!

At least two options to define graph:

Record Graph := { V : Type ; E : V → V → Type }.

Record IsGraph (V : Type) (E : V → V → Type) := { }.

Proof assistant performance (fixes)

$\text{Record Graph} := \{\, \text{V} : \text{Type}\,;\, \text{E} : \text{V} \to \text{V} \to \text{Type}\,\}.$

$\text{Record IsGraph}\ (V : \text{Type})\ (E : V \to V \to \text{Type}) := \{\,\}.$

Big difference for size of functor:

$\text{Mapping} : \text{Graph} \to \text{Graph} \to \text{Type}.$

vs.

$$\text{IsMapping} : \forall\ (V_G : \text{Type})\ (V_H : \text{Type})$$
$$(E_G : V_G \to V_G \to \text{Type})\ (E_H : V_H \to V_H \to \text{Type}),$$
$$\text{IsGraph}\ V_G\ E_G \to \text{IsGraph}\ V_H\ E_H \to \text{Type}.$$

Proof assistant performance (fixes)

- How?
  - Exceedingly careful engineering to get proofs for free

Proof assistant performance (fixes)

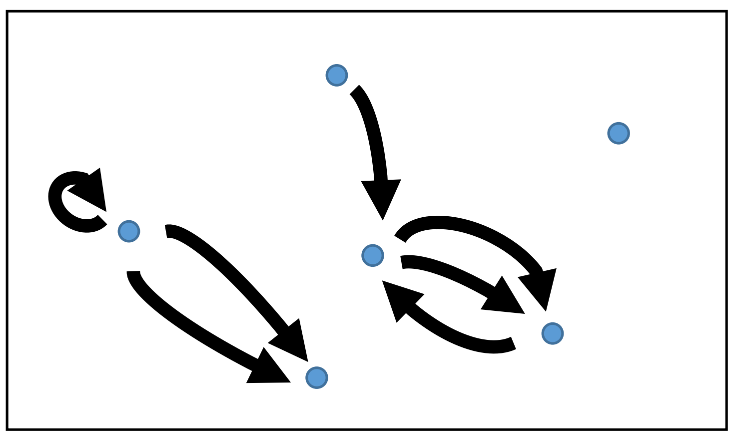
- Duality proofs for free

Proof assistant performance (fixes)

- Duality proofs for free
- Idea: One proof, two theorems

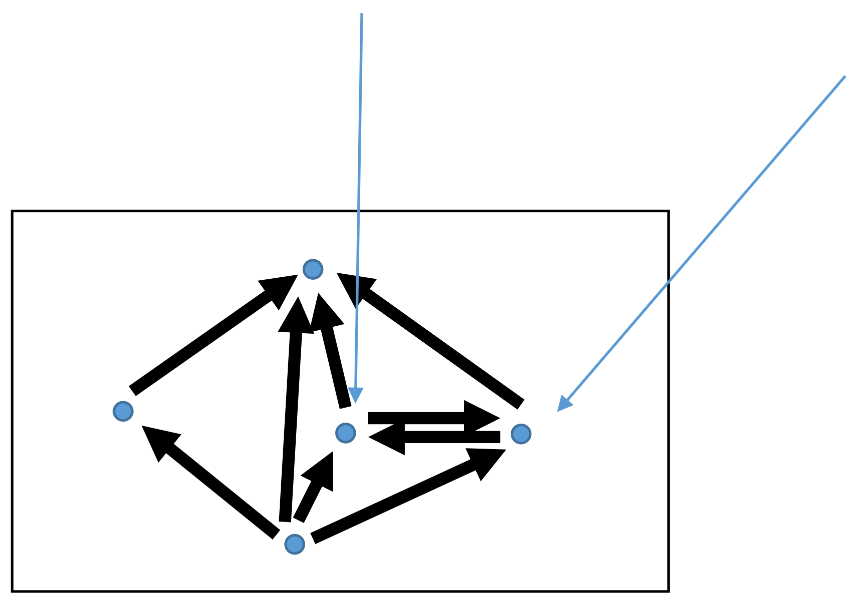
Proof assistant performance (fixes)

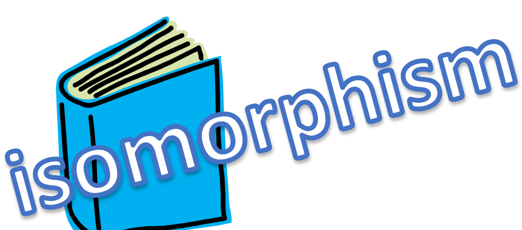
- Duality proofs for free
- Recall: A **directed graph** has:
  - a type of vertices (points)
  - for every ordered pair of vertices, a type of arrows



Proof assistant performance (fixes)

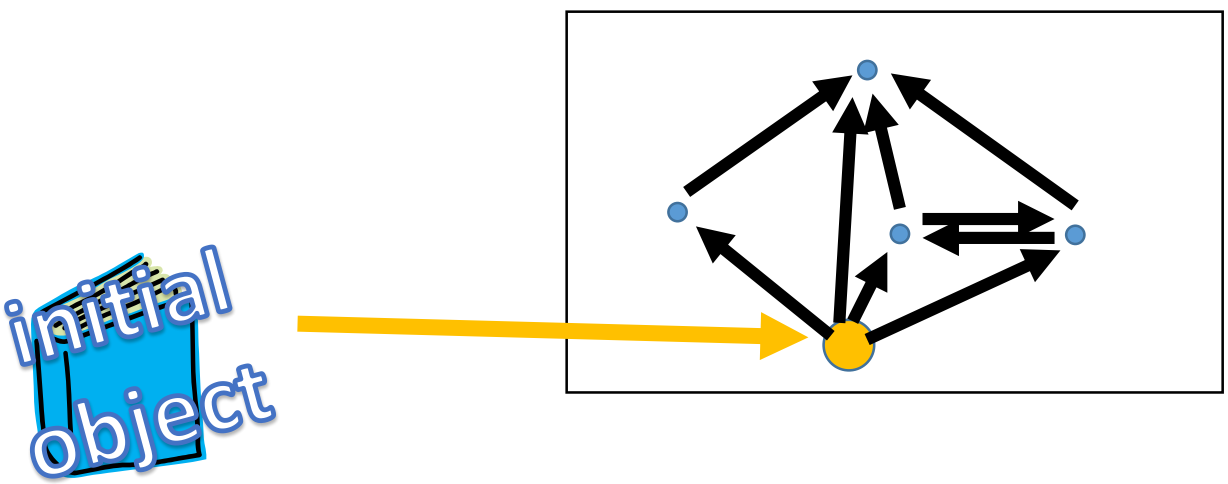
- Duality proofs for free
- Two vertices are **isomorphic** if there is exactly one edge between them in each direction





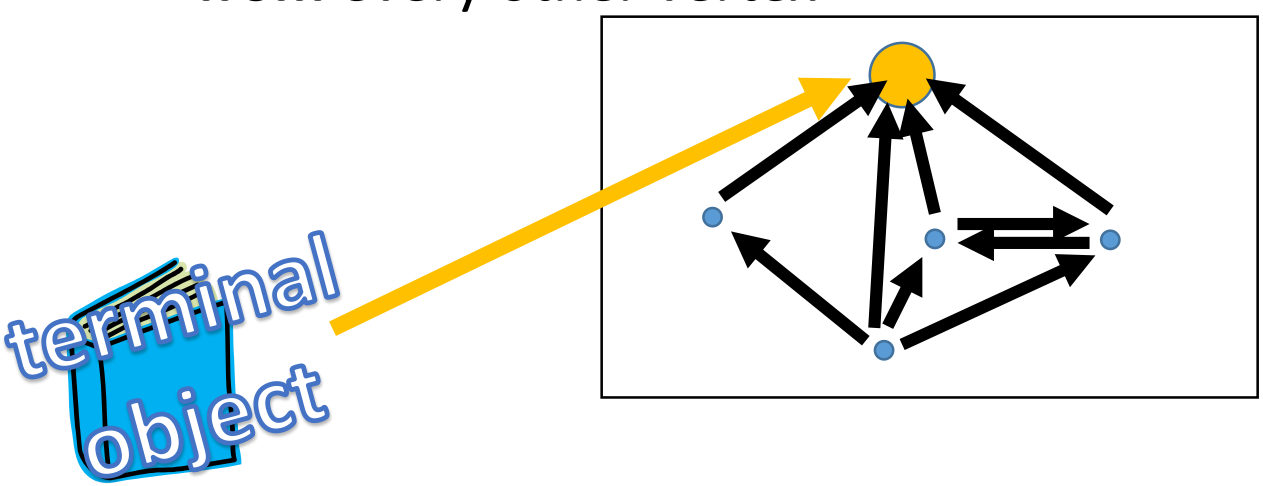
Proof assistant performance (fixes)

- Duality proofs for free
- Two vertices are **isomorphic** if there is exactly one edge between them in each direction
- An **initial (bottom) vertex** is a vertex with exactly one edge **to** every other vertex



Proof assistant performance (fixes)

- Duality proofs for free
- Two vertices are **isomorphic** if there is exactly one edge between them in each direction
- An **initial (bottom) vertex** is a vertex with exactly one edge **to** every other vertex
- A **terminal (top) vertex** is a vertex with exactly one edge **from** every other vertex



Proof assistant performance (fixes)

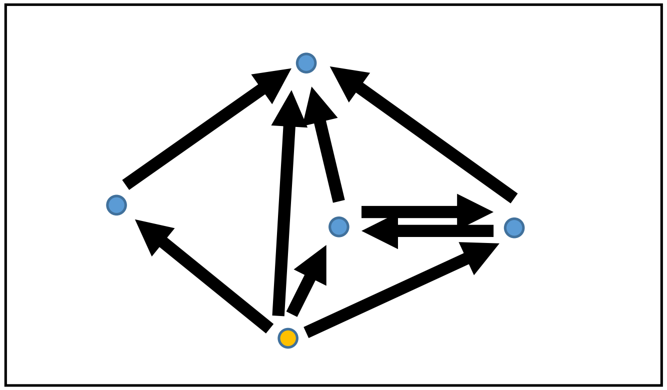
- Theorem: Initial vertices are unique

Theorem initial_unique : $\forall$ (G : Graph) (x y : G.V),

$$\text{is\_initial } x \rightarrow \text{is\_initial } y \rightarrow x \cong y$$

- Proof:

Exercise for the audience



Proof assistant performance (fixes)

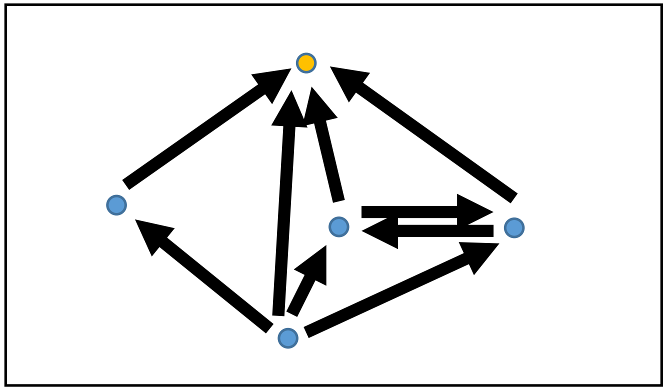
- Theorem: Terminal vertices are unique

Theorem terminal_unique : $\forall$ (G : Graph) ($x\ y$: G.V),

  is_terminal x → is_terminal y → $x \cong y$

- Proof:

  $\lambda\ G\ x\ y\ H\ H' \Rightarrow$ initial_unique $G^{\mathrm{op}}\ y\ x\ H'\ H$



Proof assistant performance (fixes)

- How?
  - Either don’t nest constructions, or don't unfold nested constructions
  - Coq only cares about unnormalized term size – “What I don't know can't hurt me”

Proof assistant performance (fixes)

- How?
  - More systematically, have good abstraction barriers

Proof assistant performance (fixes)

- How?
  - Have good abstraction barriers

Leaky abstraction barriers generally only torture programmers

Proof assistant performance (fixes)

- How?
  - Have good abstraction barriers

Leaky abstraction barriers torture Coq, too!

Proof assistant performance (fixes)

- How?
  - Have good abstraction barriers

Example: Pairing

Two ways to make use of elements of a pair:

$\text{let } (x, y) := p \text{ in } f\ x\ y$. (pattern matching)

$f\ (\text{fst } p)\ (\text{snd } p)$. (projections)

Proof assistant performance (fixes)

- How?
  - Have good abstraction barriers

Example: Pairing

Two ways to make use of elements of a pair:

$\text{let } (x, y) := p \text{ in } f\ x\ y$. (pattern matching)

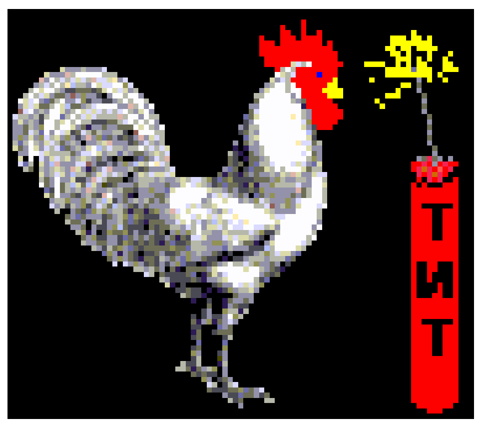
$f\ (\text{let } (x, y) := p \text{ in } x)\ (\text{let } (x, y) := p \text{ in } y)$. (projections)

These ways do not unify!

Proof assistant performance (fixes)

- How?
  - Have good abstraction barriers

Leaky abstraction barriers torture Coq, too!



Rooster Image from http://www.animationlibrary.com/animation/18342/Chicken_blows_up/

Dam image from http://www.flickr.com/photos/gammaman/7803829282/ by Eli Christman, CC by 2.0

Proof assistant performance (fixes)

- How?
  - Have good abstraction barriers

Leaky abstraction barriers torture Coq, too!

Proof assistant performance (fixes)

Concrete Example (Old Version)

```
Local Notation mor_of Y₀ Y₁ f :=
  (let η_{Y₁} := IsInitialMorphism_morphism (@HM Y₁) in
   (@center _ (IsInitialMorphism_property (@HM Y₀) _ (η_{Y₁} ∘ f))) ₁ ) (only parsing).
Lemma composition_of x y z g f : mor_of _ _ (f ∘ g) = mor_of y z f ∘ mor_of x y g.
Proof.
 simpl.
 match goal with | [ ⊢ ((@center ?A?H) ₂) ₁= _ ] ⇒ erewrite (@contr A H (center _; (_; _))) end.      8 s
 simpl; reflexivity.                                                                                   2 s
 Grab Existential Variables.
 simpl in *.
 repeat match goal with | [ ⊢ appcontext[(?x ₂) ₁ ] ] ⇒ generalize (x ₂); intro end.                   2.5 s
 rewrite ?composition_of.                                                                              0.5 s
 repeat try_associativity_quick (idtac; match goal with | [ ⊢ appcontext[?x ₁] ] ⇒ simpl rewrite x ₂ end).  3.5 s
 rewrite ?left_identity, ?right_identity, ?associativity.                                              0.3 s
 reflexivity.
Qed.                                                                                                   20 s
```

Size of goal (after first simpl): 7312 words

Size of proof term: 66 264 words

Total time in file: 39 s

Proof assistant performance (fixes)

Concrete Example (New Version)

Local Notation mor_of Y_0 Y_1 f :=
(let η_{Y_1} := IsInitialMorphism_morphism (@HM Y_1) in
IsInitialMorphism_property_morphism (@HM Y_0) _ ($\eta_{Y_1} \circ f$)) (only parsing).
Lemma composition_of x y z g f : mor_of _ _ ($f \circ g$) = mor_of y z f $\circ$ mor_of x y g.
Proof.
simpl.
erewrite IsInitialMorphism_property_morphism_unique; [reflexivity |].
rewrite ?composition_of.
repeat try_associativity_quick rewrite IsInitialMorphism_property_morphism_property.
reflexivity.
Qed.

0.08 s (was 10 s)

0.08 s (was 0.5 s)

0.5 s (was 3.5 s)

0.5 s (was 3.5 s)

Size of goal (after first simpl): 191 words (was 7312)
Size of proof term: 3 632 words (was 66 264)
Total time in file: 3 s (was 39 s)

Proof assistant performance (fixes)

Concrete Example (Old Interface)

```
Definition IsInitialMorphism_object (M : IsInitialMorphism Aφ) : D := CommaCategory.b Aφ.
Definition IsInitialMorphism_morphism (M : IsInitialMorphism Aφ) : morphism C X (U ₀ (IsInitialMorphism_object M)) := CommaCategory.f Aφ.
Definition IsInitialMorphism_property (M : IsInitialMorphism Aφ)  (Y : D) (f : morphism C X (U ₀ Y))
: Contr { m : morphism D (IsInitialMorphism_object M) Y | U ₁ m ∘ (IsInitialMorphism_morphism M) = f }.
Proof.
 (** We could just [rewrite right_identity], but we want to preserve judgemental computation rules. *)
 pose proof (@trunc_equiv' _ _ (symmetry _ _ (@CommaCategory.issig_morphism _ _ _ !X U _ _)) -2 (M (CommaCategory.Build_object !X U tt Y f))) as H'.
 simpl in H'.
 apply contr_inhabited_hprop.
 - abstract (
     apply @trunc_succ in H';
     eapply @trunc_equiv'; [ | exact H' ];
     match goal with
       | [ ⊢ appcontext[?m ∘ 𝕀] ] ⇒ simpl rewrite (right_identity _ _ _ m)
       | [ ⊢ appcontext[𝕀 ∘ ?m] ] ⇒ simpl rewrite (left_identity _ _ _ m)
     end;
     simpl; unfold IsInitialMorphism_object, IsInitialMorphism_morphism;
     let A := match goal with ⊢ Equiv ?A ?B ⇒ constr:(A) end in
     let B := match goal with ⊢ Equiv ?A ?B ⇒ constr:(B) end in
     apply (equiv_adjointify (λ x : A ⇒ x ₂) (λ x : B ⇒ (tt; x)));
     [ intro; reflexivity | intros [[]]; reflexivity ]
   ).
 - (exists ((@center _ H') ₂) ₁).
   abstract (etransitivity; [ apply ((@center _ H') ₂) ₂ | auto with morphism ]).
Defined.
```

3 s (the abstract (…) block)

1 s (the final abstract (etransitivity; …) line)

Total file time: 7 s

Proof assistant performance (fixes)

Concrete Example (New Interface)

```
Definition IsInitialMorphism_object (M : IsInitialMorphism Aφ) : D := CommaCategory.b Aφ.
Definition IsInitialMorphism_morphism (M : IsInitialMorphism Aφ) : morphism C X (U_0 (IsInitialMorphism_object M)) := CommaCategory.f Aφ.
Definition IsInitialMorphism_property_morphism (M : IsInitialMorphism Aφ)  (Y : D) (f : morphism C X (U_0 Y)) : morphism D (IsInitialMorphism_object M) Y
 := CommaCategory.h (@center _ (M (CommaCategory.Build_object !X U tt Y f))).
Definition IsInitialMorphism_property_morphism_property (M : IsInitialMorphism Aφ)  (Y : D) (f : morphism C X (U_0 Y))
: U_1 (IsInitialMorphism_property_morphism M Y f) ∘ (IsInitialMorphism_morphism M) = f
 := CommaCategory.p (@center _ (M (CommaCategory.Build_object !X U tt Y f))) @ right_identity _ _ _ _.
Definition IsInitialMorphism_property_morphism_unique (M : IsInitialMorphism Aφ)  (Y : D) (f : morphism C X (U_0 Y)) m′ (H : U_1 m′ ∘ IsInitialMorphism_morphism M = f)
: IsInitialMorphism_property_morphism M Y f = m′
:= ap (@CommaCategory.h _ _ _ _ _ _ _)
       (@contr _ (M (CommaCategory.Build_object !X U tt Y f)) (CommaCategory.Build_morphism Aφ (CommaCategory.Build_object !X U tt Y f) tt m′ (H @ (right_identity _ _ _ _) ^-1))).
Definition IsInitialMorphism_property (M : IsInitialMorphism Aφ)  (Y : D) (f : morphism C X (U_0 Y))
: Contr { m : morphism D (IsInitialMorphism_object M) Y | U_1 m ∘ (IsInitialMorphism_morphism M) = f }.
 := {| center := (IsInitialMorphism_property_morphism M Y f; IsInitialMorphism_property_morphism_property M Y f);
       contr m′ := path_sigma _ (IsInitialMorphism_property_morphism M Y f; IsInitialMorphism_property_morphism_property M Y f)
                                  m′ (@ IsInitialMorphism_property_morphism_unique M Y f m′_1 m′_2) (center _) |}.
```

0.4 s

Total file time: 7 s

Proof assistant performance (fixes)

Concrete Example 2 (Generalization)

Lemma pseudofunctor_to_cat_assoc_helper $\{x\ x_0 : C\}$ $\{x_2 :$ morphism C x x0$\}$ $\{$x1 $: C\}$

$\{x_5 :$ morphism $C\ x_0\ x_1\}$ $\{x_4 : C\}$ $\{x_7 :$ morphism $C\ x_1\ x_4\}$

$\{p\ p_0 :$ PreCategory$\}$ $\{f :$ morphism $C\ x\ x_4 \to$ Functor $p_0\ p\}$

$\{p_1\ p_2 :$ PreCategory$\}$ $\{f_0 :$ Functor $p_2\ p\}$ $\{f_1 :$ Functor $p_1\ p_2\}$ $\{f_2 :$ Functor $p_0\ p_2\}$ $\{f_3 :$ Functor $p_0\ p_1\}$ $\{f_4 :$ Functor $p_1\ p\}$

$\{x_{16} :$ morphism $(\_ \to \_)\ (f\ (x_7 \circ x_5 \circ x_2))\ (f_4 \circ f_3)$%functor$\}$

$\{x_{15} :$ morphism $(\_ \to \_)\ f_2\ (f_1 \circ f_3)$%functor$\}$ $\{H_2 :$ IsIsomorphism $x_{15}\}$

$\{x_{11} :$ morphism $(\_ \to \_)\ (f\ (x_7 \circ (x_5 \circ x_2)))\ (f_0 \circ f_2)$%functor$\}$

$\{H_1 :$ IsIsomorphism $x_{11}\}$ $\{x_9 :$ morphism $(\_ \to \_)\ f_4\ (f_0 \circ f_1)$%functor$\}$ $\{$fst_hyp $: x_7 \circ x_5 \circ x_2 = x_7 \circ (x_5 \circ x_2)\}$

(rew_hyp : $\forall\ x_3 : p_0$,

(idtoiso $(p_0 \to p)$ (ap f fst_hyp) : morphism_ _ _) $x_3 = x_{11}{}^{-1}\ x_3 \circ (f_0\ {}_1\ (x_{15}{}^{-1}\ x_3) \circ (\mathbb{I} \circ (x_9\ (f_3\ x_3) \circ x_{16}\ x_3))))$

$\{H_0' :$ IsIsomorphism $x_{16}\}$ $\{H_1' :$ IsIsomorphism $x_9\}$ $\{x_{13} : p\}$ $\{x_3 : p_0\}$ $\{x_6 : p_1\}$ $\{x_{10} : p_2\}$

$\{x_{14} :$ morphism $p\ (f_0\ x_{10})\ x_{13}\}$ $\{x_{12} :$ morphism $p_2\ (f_1\ x_6)\ x_{10}\}$ $\{x_8 :$ morphism $p_1\ (f_3\ x_3)\ x_6\}$

: existT ($\lambda\ f_5 :$ morphism $C\ x\ x_4 \Rightarrow$ morphism $p\ ((f\ f_5)\ x_3)\ x_{13}$)

$(x_7 \circ x_5 \circ x_2)$

$(x_{14} \circ (f_0\ {}_1\ x_{12} \circ x_9\ x_6) \circ (f_4\ {}_1\ x_8 \circ x_{16}\ x_3)) = (x_7 \circ (x_5 \circ x_2); x_{14} \circ (f_0\ {}_1\ (x_{12} \circ (f_1\ {}_1\ x_8 \circ x_{15}\ x_3)) \circ x_{11}\ x_3))$.

Proof.

helper_t assoc_before_commutes_tac.

assoc_fin_tac.

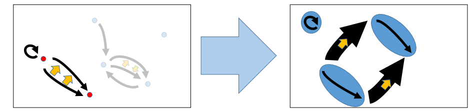
Qed.

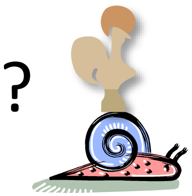
Speedup: 100x for the file, from 4m 53s to 28 s

Time spent: a few hours

Outline

- Why should we care about performance?
- What makes theorem provers (mainly Coq) slow?
  - Examples of particular slowness





- **For users (workarounds)**
  - **Arguments vs. fields and packed records**
  - **Proof by duality as proof by unification**
  - **Abstraction barriers**
  - **Proof by reflection**

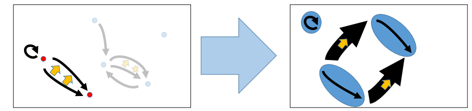
- For developers (features)
  - Primitive Projections
  - Higher inductive types
  - Universe Polymorphism
  - More judgmental rules
  - Hashconsing

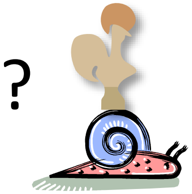
Proof assistant performance (fixes)

Better Proof Assistants

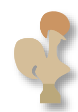
Outline

- Why should we care about performance?
- What makes theorem provers (mainly Coq) slow?
  - Examples of particular slowness

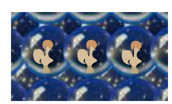




- For users (workarounds)
  - Arguments vs. fields and packed records
  - Proof by duality as proof by unification
  - Abstraction barriers
  - Proof by reflection



- **For developers (features)**
  - **Primitive projections**
  - **Universe Polymorphism**
  - **Higher inductive types**
  - **More judgmental rules**
  - **Hashconsing**



Universes image from Abell NGC2218 hst big, NASA, http://en.wikipedia.org/wiki/Abell_2218#mediaviewer/File:Abell_NGC2218_hst_big.jpg, released in Public Domain; Bubble from http://pixabay.com/en/blue-bubble-shiny-157652/, released in Public Domain CC0, combined in Photoshop by Jason Gross

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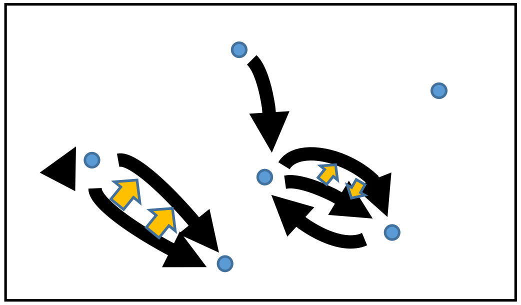
Fence image from http://www.picgifs.com/clip-art/playing-children/clip-art-playing-children-362018-689955/

Proof assistant performance (fixes)

- How?
  - Primitive projections

Proof assistant performance (fixes)

- How?
  - Primitive projections



Definition 2-Graph :=

{ V : Type &

{ 1E : V → V → Type &

$\forall\, v_1\ v_2$, 1E $v_1\ v_2$ → 1E $v_1\ v_2$ → Type }.

Definition V (G : 2-Graph) := pr_1 G .

Definition 1E (G : 2-Graph) := pr_1 (pr_2 G).

Definition 2E (G : 2-Graph) := pr_2 (pr_2 G).

Proof assistant performance (fixes)

Definition 2-Graph :=

{ V : Type &

{ 1E : V → V → Type &

∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type }.

Definition V (G : 2-Graph) := pr_1 G .

Proof assistant performance (fixes)

Definition 2-Graph :=

{ V : Type &

{ 1E : V → V → Type &

∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type }.

Definition V (G : 2-Graph) :=

@pr$_1$ Type (λ V : Type ⇒

{ 1E : V → V → Type &

∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type })

G.

Proof assistant performance (fixes)

Definition 2-Graph :=

{ V : Type &

{ 1E : V → V → Type &

∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type }.

Definition V (G : 2-Graph) := pr_1 G .

Definition 1E (G : 2-Graph) := pr_1 (pr_2 G).

Proof assistant performance (fixes)

```
Definition 1E (G : 2-Graph) :=
 @pr₁
  (@pr₁  Type (λ V : Type ⇒
                    { 1E  : V → V → Type &
                          ∀ v₁ v₂, 1E v₁ v₂ → 1E v₁ v₂ → Type })
           G →
    @pr₁  Type (λ V : Type ⇒
                    { 1E  : V → V → Type &
                          ∀ v₁ v₂, 1E v₁ v₂ → 1E v₁ v₂ → Type })
           G →
    Type)
  (λ 1E : @pr₁ Type (λ V : Type ⇒
{                1E : V → V → Type &
```

Proof assistant performance (fixes)

```
Definition 1E (G : 2-Graph) :=
 @pr_1
   (@pr_1 Type (λ V : Type ⇒
                   { 1E : V → V → Type &
                         ∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type })
          G →
    @pr_1 Type (λ V : Type ⇒
                   { 1E : V → V → Type &
                         ∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type })
          G →
    Type)
   (λ 1E : @pr_1 Type (λ V : Type ⇒
                          { 1E : V → V → Type &
                                ∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type })
                 G →
           @pr_1 Type (λ V : Type ⇒
                          { 1E : V → V → Type &
                                ∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type })
                 G →
           Type ⇒
       ∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type)
   (@pr_2 Type (λ V : Type ⇒
                   { 1E   : V → V → Type &
                          ∀ v_1 v_2, 1E v_1 v_2 → 1E v_1 v_2 → Type }
          G)
```

Proof assistant performance (fixes)

Definition 1E (G : 2-Graph) :=
@pr$_1$
(@pr$_1$ Type (λ V : Type ⇒ { 1E : V → V → Type & ∀ (v_1 : V) (v_2 : V), 1E v_1 v_2 → 1E v_1 v_2 → Type }) G →
@pr$_1$ Type (λ V : Type ⇒ { 1E : V → V → Type & ∀ (v_1 : V) (v_2 : V), 1E v_1 v_2 → 1E v_1 v_2 → Type }) G →
Type)
(λ 1E : @pr$_1$ Type (λ V : Type ⇒ { 1E : V → V → Type & ∀ (v_1 : V) (v_2 : V), 1E v_1 v_2 → 1E v_1 v_2 → Type }) G →
@pr$_1$ Type (λ V : Type ⇒ { 1E : V → V → Type & ∀ (v_1 : V) (v_2 : V), 1E v_1 v_2 → 1E v_1 v_2 → Type }) G →
Type ⇒
∀(v_1 : @pr$_1$ Type (λ V : Type ⇒ { 1E : V → V → Type & ∀ (v_1 : V) (v_2 : V), 1E v_1 v_2 → 1E v_1 v_2 → Type }) G)
(v_2 : @pr$_1$ Type (λ V : Type ⇒ { 1E : V → V → Type & ∀ (v_1 : V) (v_2 : V), 1E v_1 v_2 → 1E v_1 v_2 → Type }) G),
1E v_1 v_2 → 1E v_1 v_2 → Type)
(@pr$_2$ Type (λ V : Type ⇒ { 1E : V → V → Type & ∀ (v_1 : V) (v_2 : V), 1E v_1 v_2 → 1E v_1 v_2 → Type }) G)
:@pr$_1$ Type (λ V : Type ⇒ { 1E : V → V → Type & ∀ (v_1 : V) (v_2 : V), 1E v_1 v_2 → 1E v_1 v_2 → Type }) G →
@pr$_1$ Type (λ V : Type ⇒ { 1E : V → V → Type & ∀ (v_1 : V) (v_2 : V), 1E v_1 v_2 → 1E v_1 v_2 → Type }) G →
Type

Recall: Original was:

Definition 1E (G : 2-Graph) := pr$_1$ (pr$_2$ G).

Proof assistant performance (fixes)

- How?
  - Primitive projections
  - They eliminate the unnecessary arguments to projections, cutting down the work Coq has to do.

Proof assistant performance (fixes)

- How?
  - Don’t use setoids

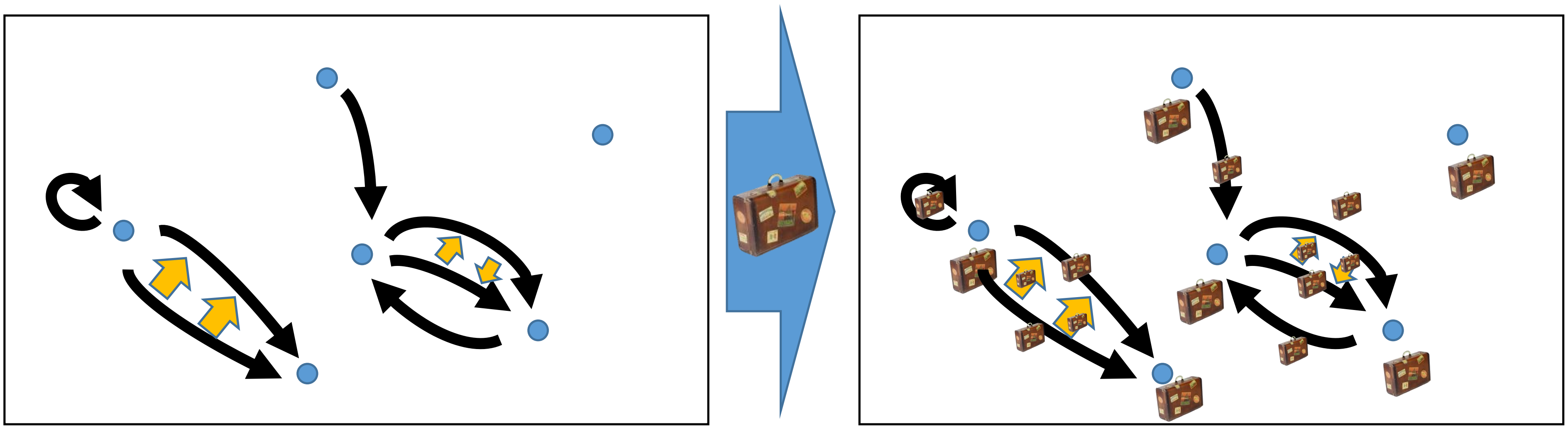
Proof assistant performance (fixes)

- How?
  - Don’t use setoids, use higher inductive types instead!

Proof assistant performance (fixes)

- How?
  - Don’t use setoids, use higher inductive types instead!

Setoids add lots of baggage to everything



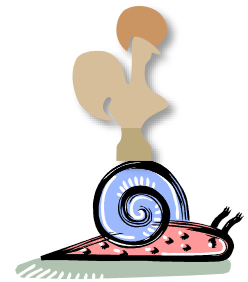
Proof assistant performance (fixes)

- How?
  - Don’t use setoids, use higher inductive types instead!

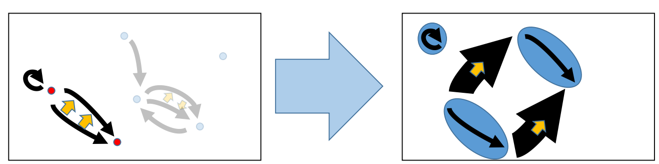
Higher inductive types (when implemented) shove the baggage into the meta-theory, where the type-checker doesn’t have to see it

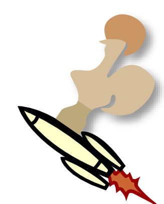
Take-away messages

- Performance matters (even in proof assistants)




- Term size matters for performance



- Performance can be improved by
  - careful engineering of developments  
  - improving the proof assistant or the metatheory  

Thank You!

The paper and presentation will be available at

http://people.csail.mit.edu/jgross/#category-coq-experience

The library is available at

https://github.com/HoTT/HoTT

subdirectory theories/categories

Questions?